#### Geometric packing and coloring problems

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joint work with



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#### References

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- F.M. de Oliveira Filho, F. Vallentin, *Fourier analysis, linear programming and densities of distance avoiding sets in* R<sup>n</sup>, (Preprint on arXiv:0808.1822v1 [math.CO], 10 pages) to appear in Journal of the European Mathematical Society

#### Motivation

In the last years I studied packing problems in compact metric spaces, e.g. the kissing number problem.

Tools: combinatorial optimization in a continuous setting:

- semidefinite optimization
- harmonic analysis



#### Questions:

- Dealing with non-compact spaces?
- From discrete point sets to sets with positive measure?

## Chromatic number of Euclidean space

 $\chi(\mathbb{R}^n)$  = minimal number of colors needed to paint all points in  $\mathbb{R}^n$  so that every two points at distance 1 receive different colors.



 $\chi_m(\mathbb{R}^n)$  = minimal number colors needed to paint all points in  $\mathbb{R}^n$  so that every two points at distance 1 receive different colors. & points receiving the same colors are Lebesgue measurable sets.

Clearly,  $\chi(\mathbb{R}^n) \leq \chi_m(\mathbb{R}^n)$ , and maybe  $\chi(\mathbb{R}^n) < \chi_m(\mathbb{R}^n)$  for some n.

## What is known?



Falconer (1981):  $\chi_m(\mathbb{R}^2) \ge 5$ 

Frankl, Wilson (1981), Larman, Rogers (1972):

 $(1 - o(1))1.2^n \le \chi(\mathbb{R}^n) \le \chi_m(\mathbb{R}^n) \le (3 + o(1))^n$ 

## Distance avoiding sets

 $A \subseteq \mathbb{R}^n$  measurable set

A avoids the distances 
$$d_1, \ldots, d_N$$
 if  
 $\forall x, y \in A : ||x - y|| \notin \{d_1, \ldots, d_N\}.$ 

A has upper density  $\overline{\delta}(A) = \limsup_{T \to \infty} \frac{\operatorname{vol}(A \cap [-T, T]^n)}{\operatorname{vol}[-T, T]^n}.$ 



#### Relation

• extreme density

 $m_{d_1,\ldots,d_N}(\mathbb{R}^n) = \sup\{\overline{\delta}(A) : A \subseteq \mathbb{R}^n \text{ avoids } d_1,\ldots,d_N\}$ 

• Upper bounds for  $m_1(\mathbb{R}^n)$  give lower bounds for  $\chi_m(\mathbb{R}^n)$ :  $m_1(\mathbb{R}^n) \cdot \chi_m(\mathbb{R}^n) \ge 1$ 

• Conjecture (Erdős):  $m_1(\mathbb{R}^2) < 0.25$ 

Best bounds (previously) known (Croft (1967), Székely (1984)):



 $0.2293 \le m_1(\mathbb{R}^2) \le 0.2790$ 

# The linear programming bound

**Main Theorem.** Suppose that there are real numbers  $z_0, z_1, \ldots, z_N$  which satisfy  $z_0 + z_1 + \cdots + z_N = 1$  and that for all  $t \ge 0$ 

 $z_0 + \Omega_n(d_1t)z_1 + \Omega_n(d_2t)z_2 + \dots + \Omega_n(d_Nt)z_N \ge 0$ holds, then  $m_{d_1,\dots,d_N}(\mathbb{R}^n) \le z_0$ .



#### Application I- one distance

#### • **LP**: $\min\{z_0 : z_0 + z_1 = 1, z_0 + \Omega_n(t)z_1 \ge 0, t \ge 0\} \ge m_1(\mathbb{R}^n)$



• Solution:  $z_0 + z_1 = 1$  and  $z_0 + \Omega_n(j_{n,1})z_1 = 0$ 

 $j_{n,1}$  — first minimum of  $\Omega_n$ 

• Corollary: 
$$\chi_m(\mathbb{R}^n) \ge 1 - \frac{1}{\Omega_n(j_{n,1})} \sim 1.165^n$$

#### Computational results

log(density) vs. dimension



# Computational results

log(colors) vs. dimension



# Application 2 - many distances

• Can be used to prove a quantitative version of

**Theorem.** (Furstenberg, Katznelson, Weiss, 1990) Let  $A \subset \mathbb{R}^2$  be a measurable set with  $\overline{\delta}(A) > 0$ . Then, there is a  $d_0$  so that for all  $d \ge d_0$  the set A does not avoid the distance d.

Their proof uses ergodic theory. Other proofs:
Bourgain (1986) — harmonic analysis
Falconer, Marstrand (1986) — geometric measure theory
Bukh (2007) — Szemerédi's regularity lemma

### Proof sketch of main theorem

 Simplifications: only avoid distance d<sub>1</sub> = 1 (saves some notation)
don't care about existence of limits (saves 10 minutes)

#### Fourier analysis

 $f,g:\mathbb{R}^n\to\mathbb{C}$  complex-valued functions

$$(f,g) = \lim_{T \to \infty} \frac{1}{\operatorname{vol}[-T,T]^n} \int_{[-T,T]^n} \frac{f(x)\overline{g(x)}dx}{\operatorname{mean value inner product}}$$

for  $u \in \mathbb{R}^n$ :  $\widehat{f}(u) = (f, e^{iu \cdot x})$  ( $\widehat{f}$  has discrete support in our case)

Fourier coefficient

 $(f,g) = \sum_{u \in \mathbb{R}^n} \widehat{f}(u)\overline{\widehat{g}(u)}$ Parseval's formula

#### Autocorrelation function

 $A \subset \mathbb{R}^n$  measurable set avoiding distance  $d_1 = 1$ 

 $\varphi(x) = (1_A, 1_{A-x})$  autocorrelation function of A  $= \lim_{T \to \infty} \frac{1}{\operatorname{vol}[-T, T]^n} \int_{[-T, T]^n} 1_A(y) 1_{A-x}(y) dy$  $= \delta(A \cap (A - x))$  $\begin{aligned} \varphi(0) &= \overline{\delta}(A) \\ \varphi(x) &= 0 \text{ if } ||x|| = 1 \end{aligned}$  $\widehat{\varphi}(u) = |\widehat{1_A}(u)|^2 \ge 0$  $= \sum |\hat{1}_A(u)|^2 e^{iu \cdot x}$  $\widehat{\varphi}(0) = \overline{\delta}(A)^2$  $u \in \mathbb{R}^n$ 

# Linear programming

The linear program in the infinite variables  $\widehat{\varphi}(u)$ 

 $\sup \left\{ \begin{aligned} & \frac{\widehat{\varphi}(0)}{\sum_{u \in \mathbb{R}^n} \widehat{\varphi}(u)} & : \quad \widehat{\varphi}(u) \ge 0, \quad u \in \mathbb{R}^n, \\ & & \sum_{u \in \mathbb{R}^n} \widehat{\varphi}(u) e^{iu \cdot x} = 0, \ \|x\| = 1 \end{aligned} \right\} \\ \text{gives an upper bound on } m_1(\mathbb{R}^n). \end{aligned}$ 



Taking spherical averages simplifies its computation tremendously (and explains the appearance of  $\Omega_n$ ).

Transform  $\varphi$  into a radial function by taking averages over the unit sphere For  $x \in \mathbb{R}^n$  define

Duality

Symmetrization gives linear program in infinite variables  $\alpha(t), t \in \mathbb{R}_{\geq 0}$ .

$$\sup \left\{ \frac{\alpha(0)}{\sum_{t \ge 0} \alpha(t)} : \alpha(t) \ge 0, \quad t \ge 0, \\ \sum_{t \ge 0} \alpha(t) \Omega_n(t) = 0 \right\}$$

Taking the dual linear program and showing weak duality proves the theorem.