

Strategic Games: Social Optima and Nash Equilibria

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Part I

- Strategic games.
- Nash equilibrium.
- Social optimum.
- Price of anarchy.
- Price of stability.

Strategic Games

Strategic game for $|N| \geq 2$ players:

$$G := (N, \{S_i\}_{i \in N}, \{p_i\}_{i \in N}).$$

For each player i

- (possibly infinite) set S_i of **strategies**,
- **payoff function** $p_i : S_1 \times \dots \times S_n \rightarrow \mathbb{R}$.

Basic assumptions

- Players choose their strategies **simultaneously**,
- each player is **rational**: his objective is to maximize his payoff,
- players have **common knowledge** of the game and of each others' rationality.

Three Examples (1)

The Battle of the Sexes

	<i>F</i>	<i>B</i>
<i>F</i>	2, 1	0, 0
<i>B</i>	0, 0	1, 2

Matching Pennies

	<i>H</i>	<i>T</i>
<i>H</i>	1, -1	-1, 1
<i>T</i>	-1, 1	1, -1

Prisoner's Dilemma

	<i>C</i>	<i>D</i>
<i>C</i>	2, 2	0, 3
<i>D</i>	3, 0	1, 1

Main Concepts

- **Notation:** $s_i, s'_i \in S_i$,
 $s, s', (s_i, s_{-i}) \in S_1 \times \dots \times S_n$.

- s is a **Nash equilibrium** if

$$\forall i \in \{1, \dots, n\} \quad \forall s'_i \in S_i \quad p_i(s_i, s_{-i}) \geq p_i(s'_i, s_{-i}).$$

- **Social welfare** of s :

$$SW(s) := \sum_{j=1}^n p_j(s).$$

- s is a **social optimum** if $SW(s)$ is maximal.

Intuitions

- **Nash equilibrium:**
Every player is 'happy'
(played his **best response**).
- **Social optimum:**
The desired state of affairs for the society.
- **Main problem:**
Social optima may not be Nash equilibria.

Three Examples (2)

The Battle of the Sexes: Two Nash equilibria.

	<i>F</i>	<i>B</i>
<i>F</i>	2,1	0,0
<i>B</i>	0,0	1,2

Matching Pennies: No Nash equilibrium.

	<i>H</i>	<i>T</i>
<i>H</i>	1, -1	-1, 1
<i>T</i>	-1, 1	1, -1

Prisoner's Dilemma: One Nash equilibrium.

	<i>C</i>	<i>D</i>
<i>C</i>	2,2	0,3
<i>D</i>	3,0	1,1

Price of Anarchy and of Stability

- **Price of Anarchy** (Koutsoupias, Papadimitriou, 1999):

$$\frac{\text{SW of social optimum}}{\text{SW of the worst Nash equilibrium}}$$

- **Price of Stability** (Schulz, Moses, 2003):

$$\frac{\text{SW of social optimum}}{\text{SW of the best Nash equilibrium}}$$

Examples

A 3×3 game

	<i>L</i>	<i>M</i>	<i>R</i>
<i>T</i>	2,2	4,1	1,0
<i>C</i>	1,4	3,3	1,0
<i>B</i>	0,1	0,1	1,1

$$\text{PoA} = \frac{6}{2} = 3.$$

$$\text{PoS} = \frac{6}{4} = 1.5.$$

Prisoner's Dilemma

	<i>C</i>	<i>D</i>
<i>C</i>	2,2	0,3
<i>D</i>	3,0	1,1

$$\text{PoA} = \text{PoA} = 2.$$

Cournot Competition (1838)

- One infinitely divisible product (oil),
- n companies decide **simultaneously** how much to produce,
- price is decreasing in total output.

Each $S_i = \mathbb{R}_+$,

$$p_i(s) := s_i \left(a - b \sum_{j=1}^n s_j \right) - cs_i$$

for some a, b, c , where $a > c$ and $b > 0$.

The price of the product: $a - b \sum_{j=1}^n s_j$.

The production cost: cs_i .

Cournot Competition (ctd)

- $p_i(s) := s_i \left(a - b \sum_{j=1}^n s_j \right) - cs_i.$

- **Unique Nash equilibrium:**

s , with each $s_i = \frac{a-c}{b(n+1)}.$

$$SW(s) = \frac{(a-c)^2}{b} \cdot \frac{n}{(n+1)^2}.$$

- **Social optimum:**

when $\sum_{j=1}^n s_j = \frac{a-c}{2b}.$

$$SW(s) = \frac{(a-c)^2}{4b}.$$

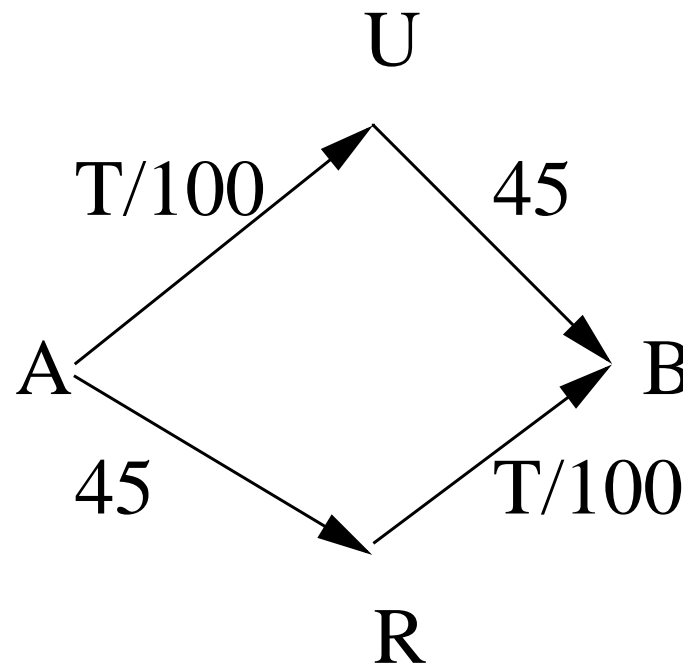
- **Note**

$$\text{PoA (= PoS)} = \frac{(n+1)^2}{4n}.$$

Congestion Games: Example

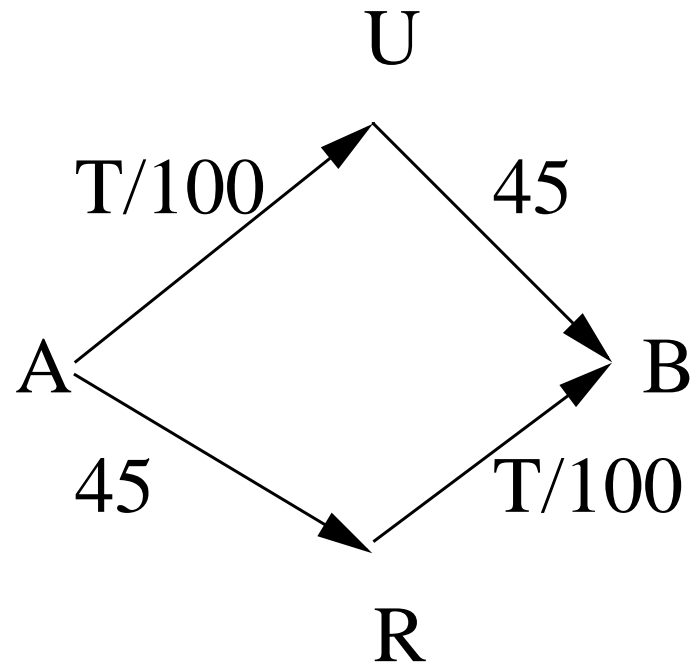
Assumptions:

- 4000 **drivers** drive from A to B.
- Each driver has 2 possibilities (**strategies**).



Problem: Find a Nash equilibrium (T = number of drivers).

Nash Equilibrium

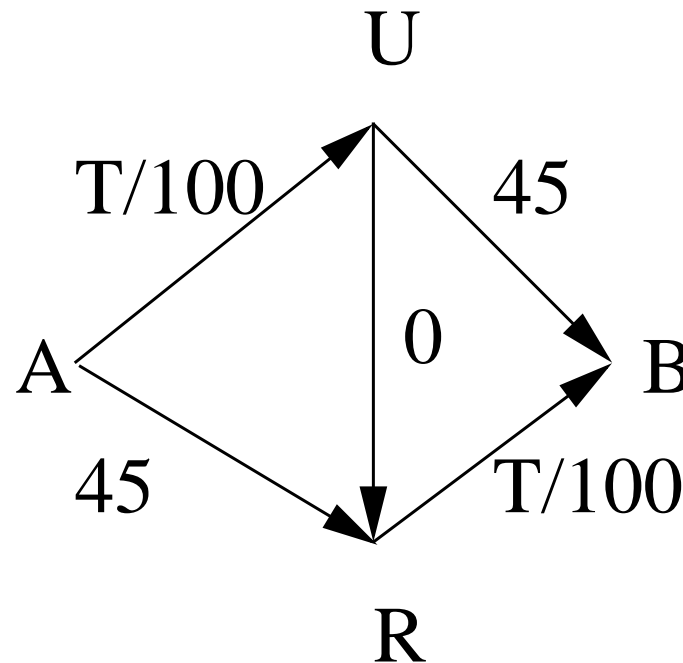


Answer: 2000/2000.

Travel time: $2000/100 + 45 = 45 + 2000/100 = 65$.

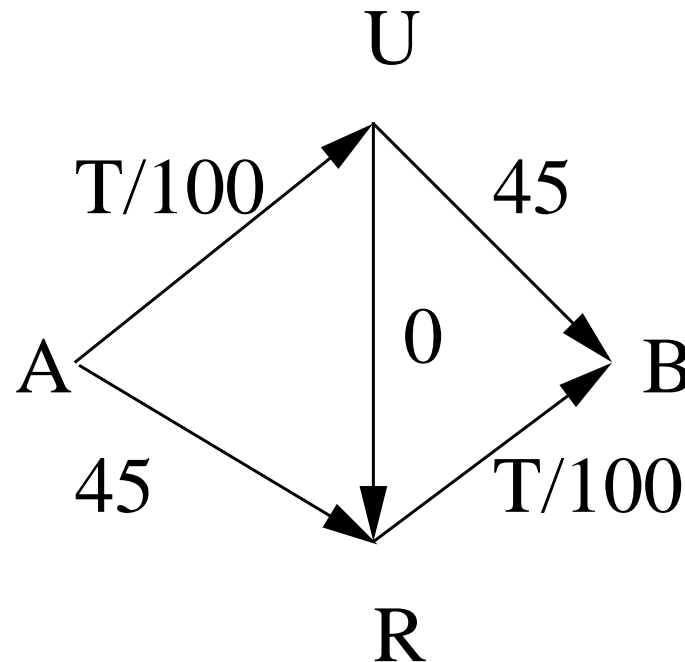
Braess Paradox

- Add a fast road from U to R.
- Each driver has now 3 possibilities (**strategies**):
A - U - B,
A - R - B,
A - U - R - B.



Problem: Find a Nash equilibrium.

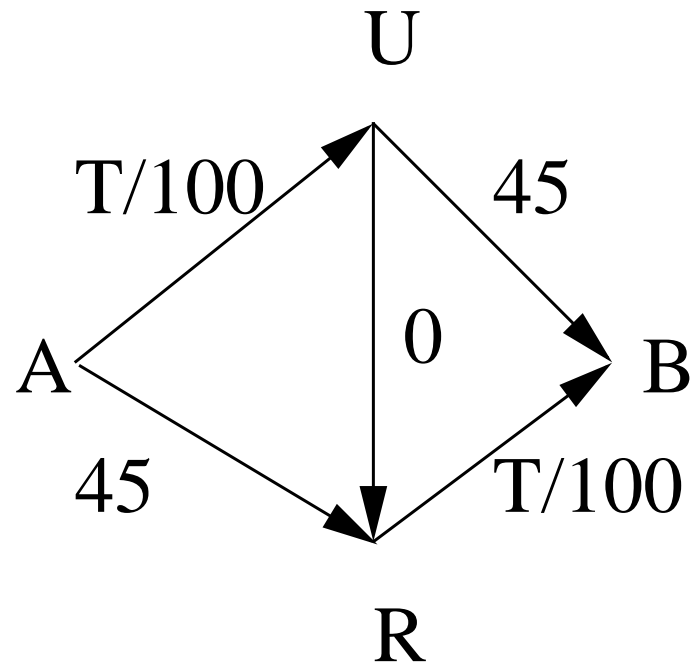
Nash Equilibrium



Answer: Each driver will choose the road A - U - R - B.

Why?: The road A - U - R - B is **always** a best response.

Bad News



- **Travel time:** $4000/100 + 4000/100 = 80!$
- **PoA (and PoS)** went up from 1 to $80/65$.

Does it Happen?

From Wikipedia ('Braess Paradox'):

- In **Seoul, South Korea**, a speeding-up in traffic around the city was seen when a motorway was removed as part of the Cheonggyecheon restoration project.
- In **Stuttgart, Germany** after investments into the road network in 1969, the traffic situation did not improve until a section of newly-built road was closed for traffic again.
- In 1990 the closing of 42nd street in **New York City** reduced the amount of congestion in the area.
- In 2008 Youn, Gastner and Jeong demonstrated specific routes in **Boston, New York City** and **London** where this might actually occur and pointed out roads that could be closed to reduce predicted travel times.

Part II

- Altruistic games.
- Selfishness level.
(Based on
Selfishness level of strategic games,
K.R. Apt and G. Schäfer)

Altruistic Games

- Given $G := (N, \{S_i\}_{i \in N}, \{p_i\}_{i \in N})$ and $\alpha \geq 0$.
- $G(\alpha) := (N, \{S_i\}_{i \in N}, \{r_i\}_{i \in N})$, where

$$r_i(s) := p_i(s) + \alpha SW(s).$$

- When $\alpha > 0$ the payoff of each player in $G(\alpha)$ depends on the social welfare of the players.
- $G(\alpha)$ is an **altruistic version** of G .

Selfishness Level (1)

- G is α -selfish if a Nash equilibrium of $G(\alpha)$ is a social optimum of $G(\alpha)$.
- If for no $\alpha \geq 0$, G is α -selfish, then its selfishness level is ∞ .
- Suppose G is finite.
If for some $\alpha \geq 0$, G is α -selfish, then

$$\min_{\alpha \in \mathbb{R}_+} (G \text{ is } \alpha\text{-selfish})$$

is the selfishness level of G .

Selfishness Level (2)

- Suppose G is **infinite**.

- If for some $\alpha \geq 0$, G is α -selfish **and**

$$\min_{\alpha \in \mathbb{R}_+} (G \text{ is } \alpha\text{-selfish})$$

exists, then it is the **selfishness level** of G .

- Otherwise the **selfishness level** of G is undefined.

Three Examples (1)

The Battle of the Sexes

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Matching Pennies

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<i>D</i>	3, 0	1, 1

Three Examples (2)

The Battle of the Sexes: selfishness level is 0.

	F	B
F	2, 1	0, 0
B	0, 0	1, 2

Matching Pennies: selfishness level is ∞ .

	H	T
H	1, -1	-1, 1
T	-1, 1	1, -1

Prisoner's Dilemma: selfishness level is 1.

	C	D
C	2, 2	0, 3
D	3, 0	1, 1

	C	D
C	6, 6	3, 6
D	6, 3	3, 3

Selfishness Level vs Price of Stability

- **Note**

Selfishness level of a finite game is 0 iff price of stability is 1.

- **Theorem**

For every finite $\alpha > 0$ and $\beta > 1$ there is a finite game with selfishness level α and price of stability β .

Example: Traveler's Dilemma

Two players, $S_i = \{2, \dots, 100\}$,

$$p_i(s) := \begin{cases} s_i & \text{if } s_i = s_{-i} \\ s_i + 2 & \text{if } s_i < s_{-i} \\ s_{-i} - 2 & \text{otherwise.} \end{cases}$$

Problem: Find a Nash equilibrium.

Proposition Selfishness level is $\frac{1}{2}$.

Cournot Competition

- **Note** Price of anarchy (and of stability) converges with n to ∞ .
- **Proposition** For each $n > 1$ the selfishness level is ∞ .

Tragedy of the Commons

- Contiguous common resource (shared bandwidth),
- Each $S_i = [0, 1]$,
- s_i : chosen fraction of the common resource
- payoff function:

$$p_i(s) := \begin{cases} s_i(1 - \sum_{j=1}^n s_j) & \text{if } \sum_{j=1}^n s_j \leq 1 \\ 0 & \text{otherwise} \end{cases}$$

- **Intuition**: the payoff degrades when the resource is overused.
- **Proposition** For each $n > 1$ the selfishness level is ∞ .

Congestion Games

Congestion game: $G = (N, E, \{S_i\}_{i \in N}, \{d_e\}_{e \in E})$,
where

- E is a finite set of facilities,
- $S_i \subseteq 2^E$ is the set of facility subsets available to player i ,
- $d_e \in \mathbb{N}$ is the delay function for facility $e \in E$.
- Let $x_e(s)$ be the number of players using facility e in s .
- The goal of a player is to minimize his individual cost $c_i(s) := \sum_{e \in s_i} d_e(x_e(s))$.
- Social cost:
 $SC(s) = \sum_{i=1}^n c_i(s)$.

Linear Congestion Games

- **Linear congestion game**: each delay function is of the form $d_e(x) = a_e x + b_e$, where $a_e, b_e \in \mathbb{N}$.
- Let L be the maximum number of facilities that any player can choose:
$$L := \max_{i \in N, s_i \in S_i} |s_i|.$$
- $\Delta_{\max} := \max_{e \in E} (a_e + b_e),$
 $\Delta_{\min} := \min_{e \in E} (a_e + b_e).$

Proposition Selfishness level of a linear congestion game is $\leq \frac{1}{2}(L \cdot \Delta_{\max} - \Delta_{\min} - 1).$

Note This bound does not depend on the number of players.

Take Home Message

- Price of anarchy and price of stability are **descriptive** concepts.
- Selfishness level is a **normative** concept.

Some Quotations

- Dalai Lama:

The intelligent way to be selfish is to work for the welfare of others.

Microeconomics: Behavior, Institutions, and Evolution, S. Bowles '04.

- An excellent way to promote cooperation in a society is to teach people to care about the welfare of others.

The Evolution of Cooperation, R. Axelrod, '84.

THANK YOU

Dziękuję za uwagę