

# Introduction to Modern Cryptography, Exercise # 8

University of Amsterdam, Master of Logic

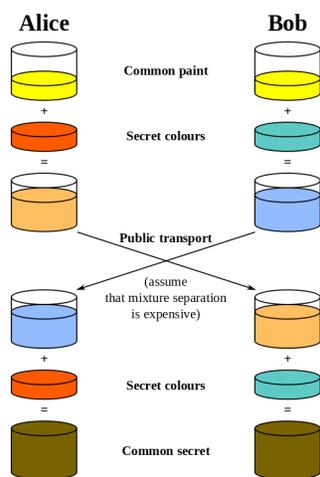
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1 November 2011

(to be handed in by Tuesday, 8 November 2011, 9:00)

1. **Square-And-Multiply, Efficient Modular Exponentiation:** Exercise B.3 in [KL]. Argue why your algorithm is efficient. **Corrected hint:** Let  $y = [a^b \bmod N]$  denote the answer. Use auxiliary variables  $x$  (initialized to  $a$ ) and  $t$  (initialized to 1), and maintain the invariant  $t \cdot x^b = y \bmod N$  while decreasing  $b$  and squaring  $x$ . The algorithm terminates when  $b = 0$  and  $t$  is equal to the answer.
2. **Interactive Secure Encryption:** Exercise 9.1 in [KL]
3. **Man-In-The-Middle Attacks:** Exercise 9.2 in [KL]
4. **Key Exchange with Bit Strings:** Exercise 9.3 in [KL]
5. **CDH and DDH:**
  - (a) Give an example of a (not necessarily multiplicative) group  $\mathcal{G}$  relative to which the CDH-Problem is easy.
  - (b) Prove formally that the hardness of the CDH problem relative to a group  $\mathcal{G}$  implies the hardness of the discrete logarithm problem relative to  $\mathcal{G}$ . (Exercise 7.15 in [KL])
  - (c) Prove formally that the hardness of the DDH problem relative to a group  $\mathcal{G}$  implies the hardness of the CDH problem relative to  $\mathcal{G}$ . (Exercise 7.16 in [KL])



Diffie-Hellman Key Exchange Using Buckets of Paint

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