

0	00	000	0000
			0001
	001	0010	
		0011	
	01	010	0100
			0101
011	0110		
	0111		
1	10	100	1000
			1001
	101	1010	
		1011	
	11	110	1100
			1101
111	1110		
	1111		

The total symbol code budget

Figure 5.1. The symbol coding budget. The 'cost'  $2^{-l}$  of each codeword (with length  $l$ ) is indicated by the size of the box it is written in. The total budget available when making a uniquely decodeable code is 1. You can think of this diagram as showing a *codeword supermarket*, with the codewords arranged in aisles by their length, and the cost of each codeword indicated by the size of its box on the shelf. If the cost of the codewords that you take exceeds the budget then your code will not be uniquely decodeable.

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$a_i$	$c(a_i)$	$l_i$
a	1000	4
b	0100	4
c	0010	4
d	0001	4

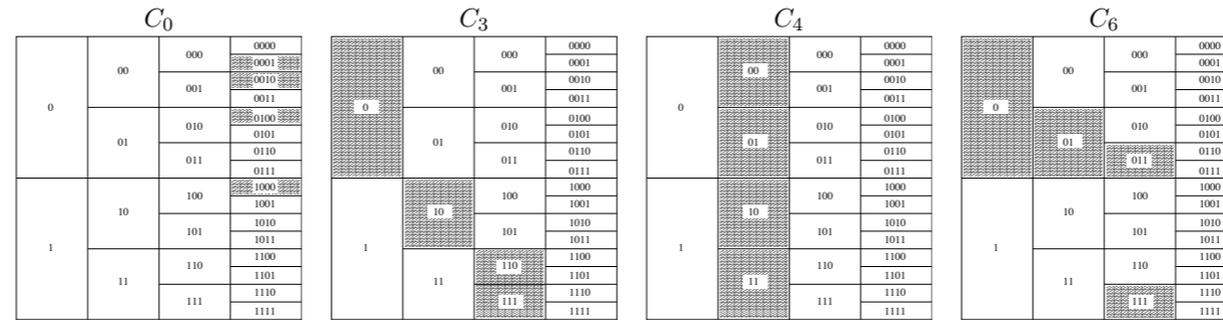
$C_0$ :

$C_3$ :				
$a_i$	$c(a_i)$	$p_i$	$h(p_i)$	$l_i$
a	0	1/2	1.0	1
b	10	1/4	2.0	2
c	110	1/8	3.0	3
d	111	1/8	3.0	3

$C_4$	$C_5$
a	00 0
b	01 1
c	10 00
d	11 11



$C_6$ :				
$a_i$	$c(a_i)$	$p_i$	$h(p_i)$	$l_i$
a	0	1/2	1.0	1
b	01	1/4	2.0	2
c	011	1/8	3.0	3
d	111	1/8	3.0	3



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$a_i$	$p_i$	$\log_2 \frac{1}{p_i}$	$l_i$	$c(a_i)$
a	0.0575	4.1	4	0000
b	0.0128	6.3	6	001000
c	0.0263	5.2	5	00101
d	0.0285	5.1	5	10000
e	0.0913	3.5	4	1100
f	0.0173	5.9	6	111000
g	0.0133	6.2	6	001001
h	0.0313	5.0	5	10001
i	0.0599	4.1	4	1001
j	0.0006	10.7	10	1101000000
k	0.0084	6.9	7	1010000
l	0.0335	4.9	5	11101
m	0.0235	5.4	6	110101
n	0.0596	4.1	4	0001
o	0.0689	3.9	4	1011
p	0.0192	5.7	6	111001
q	0.0008	10.3	9	110100001
r	0.0508	4.3	5	11011
s	0.0567	4.1	4	0011
t	0.0706	3.8	4	1111
u	0.0334	4.9	5	10101
v	0.0069	7.2	8	11010001
w	0.0119	6.4	7	1101001
x	0.0073	7.1	7	1010001
y	0.0164	5.9	6	101001
z	0.0007	10.4	10	1101000001
-	0.1928	2.4	2	01

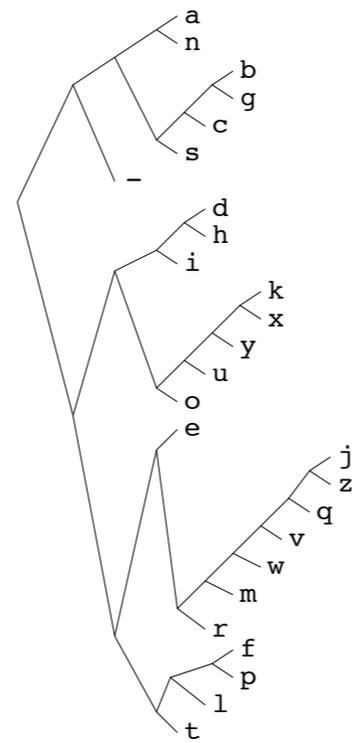


Figure 5.6. Huffman code for the English language ensemble (monogram statistics).

Figure 5.6. This code has an expected length of 4.15 bits; the entropy of the ensemble is 4.11 bits. Observe the disparities between the assigned codelengths and the ideal codelengths  $\log_2 \frac{1}{p_i}$ .

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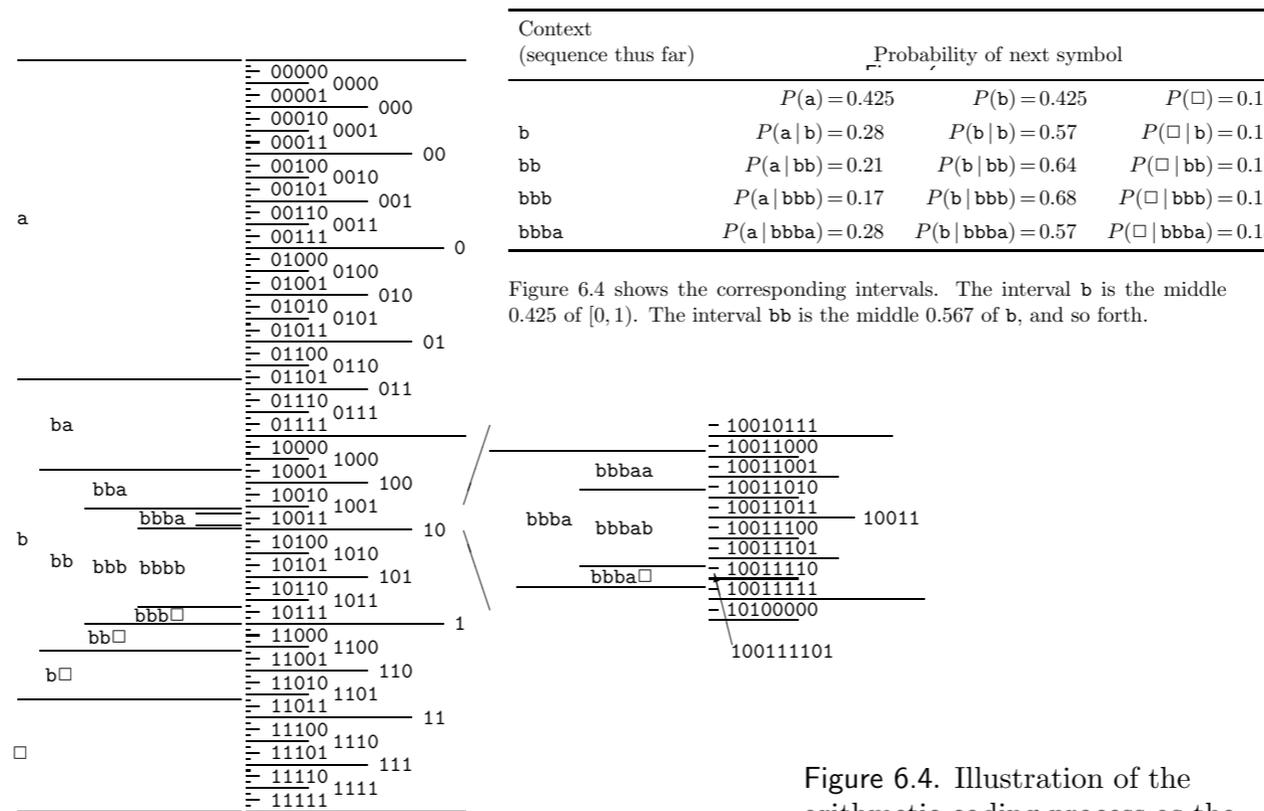


Figure 6.4. Illustration of the arithmetic coding process as the sequence **bbba** $\square$  is transmitted.

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