

Information Theory Exercise Sheet #7

(Channel Capacities)

University of Amsterdam, Master of Logic, Fall 2015

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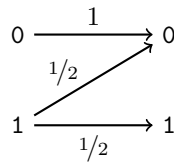
Out: Wednesday, 9 December 2015
(due: Friday, 18 December 2015, 9:00)

1. **Channel Capacities (Cover and Thomas, Ex. 7.34)** Find the capacity and optimal input distribution of the following channels:

- (a) The Z-channel with transition matrix

$$P_{Y|X} = \begin{bmatrix} 1 & 0 \\ 1/2 & 1/2 \end{bmatrix}$$

and channel diagram



- (b) The ternary channel with transition matrix

$$P_{Y|X} = \begin{bmatrix} 2/3 & 1/3 & 0 \\ 1/3 & 1/3 & 1/3 \\ 0 & 1/3 & 2/3 \end{bmatrix}$$

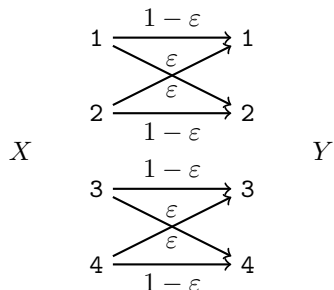
2. **Toggle channel (Cover and Thomas, Ex. 7.28)** Suppose we are given two (possibly different) channels with capacities C_1 and C_2 . We then define a new “union channel” by allowing the transmitter at each time to choose between sending a signal through either channel 1 or channel 2. We assume that the output alphabets of the two component channels are distinct and do not intersect.

- (a) Letting C stand for the capacity of the union channel, show that

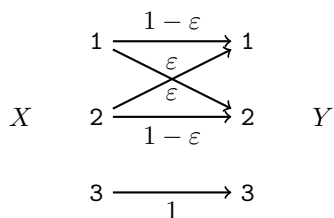
$$2^C = 2^{C_1} + 2^{C_2}.$$

Use this result to calculate the capacity of the following channels:

(b) Two parallel BSCs:



(c) A BSC and a single symbol:



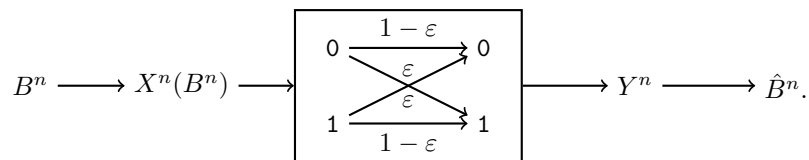
3. **Noise alphabets (Cover and Thomas, Ex. 7.30)** Consider an additive-noise channel where $\mathcal{X} = \{0, 1, 2, 3\}$ and $Y = X + Z$ with addition over the integers \mathbb{Z} . Z is uniformly distributed over three distinct integer values $\mathcal{Z} = \{z_1, z_2, z_3\}$.

- What is the maximum capacity over all choices of the \mathcal{Z} alphabet? Give distinct integer values z_1, z_2, z_3 and a distribution on \mathcal{X} achieving this.
- What is the minimum capacity over all choices of the \mathcal{Z} alphabet? Give distinct integer values z_1, z_2, z_3 and a distribution on \mathcal{X} achieving this.

4. **Source and channel (Cover and Thomas, Ex. 7.31)** We wish to encode a sequence $B_1, B_2, B_3, \dots, B_n$ of i.i.d. Bernoulli(θ) variables,

$$P_B(1) = 1 - P_B(0) = \theta,$$

for transmission over a binary symmetric channel with flip probability ε ,



Find conditions on (θ, ε) under which the error probability $\Pr[\hat{B}^n \neq B^n]$ can be made to go to zero as $n \rightarrow \infty$.

5. **Another Kind of Entropy.** In this exercise we consider a different entropy notion. Let X and Y be random variables with joint probability distribution P_{XY} . The *guessing probability* and the *min-entropy* of X are respectively defined as

$$\text{Guess}(X) := \max_x P_X(x) \quad \text{and} \quad H_{\min}(X) := -\log \text{Guess}(X).$$

The *conditional guessing probability* and the *conditional min-entropy* of X are respectively defined as

$$\text{Guess}(X|Y) := \sum_y P_Y(y) \text{Guess}(X|Y = y)$$

and

$$H_{\min}(X|Y) := -\log \text{Guess}(X|Y).$$

- (a) If X has no uncertainty (i.e. $H(X) = 0$), what is $H_{\min}(X)$?
- (b) If X is uniformly distributed over \mathcal{X} , what is $H_{\min}(X)$?
- (c) Prove that $H_{\min}(XY) \geq H_{\min}(X)$.
- (d) Prove that $H_{\min}(X) \geq H_{\min}(X|Y)$.
- (e) Prove that $H_{\min}(X|Y) \geq H_{\min}(XY) - \log |\mathcal{Y}|$.

Homework is exercises 1, 2, 3, and 4.