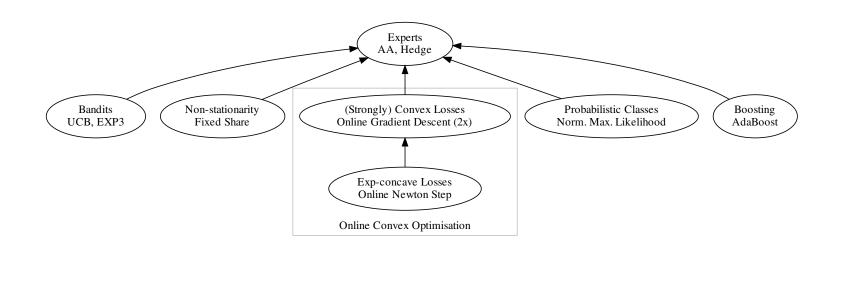
Machine Learning Theory. Lecture 10. Wouter M. Koolen

Recap:

• Mix loss (now with non-uniform regret bounds)

Non-stationary environments

• Switching (Fixed Share algorithm)



Mix loss prediction (w. non-uniform regret objective)

For t = 1, 2, ...

- 1. Play $w_t \in \triangle_K$.
- 2. See $\ell_t \in \mathbb{R}^K$.
- 3. Incur mix loss $\hat{\ell}_t \coloneqq -\ln\left(\sum_{k=1}^K w_t^k e^{-\ell_t^k}\right)$.

Definition 1. The regret w.r.t. expert $k \in [K]$ after $T \ge 0$ rounds is

$$R_T^k \coloneqq \sum_{t=1}^T \left(\hat{\ell}_t - \ell_t^k\right).$$

The Aggregating Algorithm with prior

Definition 2. *The Aggregating Algorithm with* prior $\pi \in \triangle_K$ *plays*

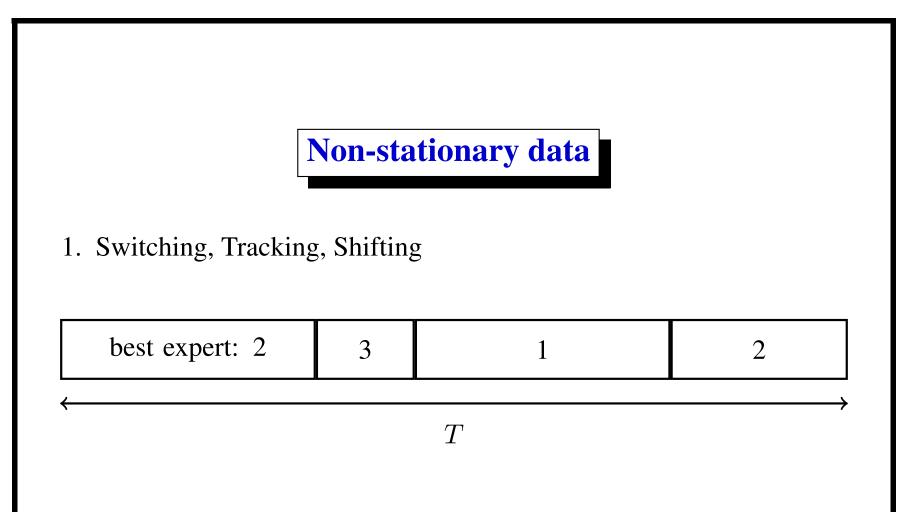
$$w_t^k = \frac{\pi^k e^{-\sum_{s=1}^{t-1} \ell_s^k}}{\sum_{j=1}^K \pi^j e^{-\sum_{s=1}^{t-1} \ell_s^j}}$$
(AA- π)

(so
$$w_1 = \pi$$
 and $w_{t+1}^k = \frac{w_t^k e^{-\ell_t^k}}{\sum_{j=1}^K w_t^j e^{-\ell_t^j}}$)

Theorem 3. The regret of $AA - \pi$ w.r.t. expert $k \in [K]$ satisfies

$$R_T^k \leq -\ln \pi^k$$

Proof: part of Homework 10.



What if no single expert is good for all T rounds, but data can be split into 4 blocks with different best experts?

Non-stationary data

So far we have been looking at regret compared to a *fixed* expert/action.

$$R_T = \max_{k \in [K]} \sum_{t=1}^T \left(\hat{\ell}_t - \ell_t^k \right).$$

But what if we do not expect a single expert to be good for all data?

Definition 4. Let $\xi \in [K]^T$ be a sequence of experts. The sequence regret w.r.t. ξ is defined by

$$R_{\xi} \coloneqq \sum_{t=1}^{T} \left(\hat{\ell}_t - \ell_t^{\xi_t} \right)$$

Question: Can we keep the sequence regret small w.r.t. *every sequence*? Or, at least, w.r.t. many interesting sequences?

Fixed Share (defined by reduction)

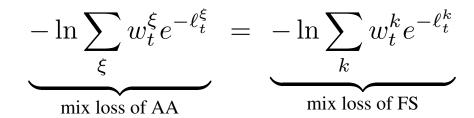
Starting with K experts, create an "explosion" $[K]^T$ of *expert sequences*. Fix *switching rate* $\alpha \in [0, 1]$. Run the AA with prior $\pi \in \Delta_{[K]^T}$ and losses $\ell_t \in (-\infty, \infty]^{[K]^T}$ defined by

$$\pi^{\xi} \coloneqq \frac{1}{K} \prod_{t=2}^{T} \begin{cases} 1 - \alpha & \text{if } \xi_{t-1} = \xi_t \\ \frac{\alpha}{K-1} & \text{if } \xi_{t-1} \neq \xi_t \end{cases}$$
$$\ell_t^{\xi} \coloneqq \ell_t^{\xi_t}$$

Based on the AA predictions $w_t^{\xi} \in \triangle_{[K]^T}$, Fixed Share plays

$$w_t^k \coloneqq \sum_{\xi \in [K]^T : \xi_t = k} w_t^{\xi}$$
(FS)

Crucial equality:



So we can directly apply AA regret bound to obtain FS regret bound.

Fixed Share: Regret Bound

Application of the AA regret bound, (Theorem 3) gives

Theorem 5. Fixed Share ensures that the regret on each sequence ξ with $B \leq T$ contiguous blocks is at most

$$R_{\xi} \leq -\ln \pi^{\xi}$$

$$= \underbrace{\ln K + (B-1)\ln(K-1)}_{expert \ labelling \ cost} \underbrace{-(B-1)\ln\alpha - (T-B)\ln(1-\alpha)}_{switching \ location \ cost}.$$

Corollary 6. If we know the number of blocks B in advance, we can optimise the bound by setting $\alpha = \frac{B-1}{T-1}$ to find

$$R_{\xi} \leq \ln K + (B-1)\ln(K-1) + (T-1)H\left(\frac{B-1}{T-1}\right),$$

where $H(\alpha) = -\alpha \ln \alpha - (1 - \alpha) \ln(1 - \alpha)$ is the binary entropy.

Fixed Share: Computation Collapses

Seems we need to maintain exponentially many weights. But prior π^{ξ} is *Markov*

Theorem 7. The weights of Fixed Share with switching rate α can be computed incrementally in $\mathcal{O}(K)$ time per round (same as AA) as

$$w_{t+1}^{k} = (1-\beta) \frac{w_{t}^{k} e^{-\ell_{t}^{k}}}{\sum_{j} w_{t}^{j} e^{-\ell_{t}^{j}}} + \frac{\beta}{K}$$

where $\beta = \alpha \frac{K}{K-1}$.

Proof: part of Homework 10.

We see that the Fixed Share update is a weighted combination of the incremental AA update and the uniform prior.

Conclusion

Technique for adapting to changing environment

• Fixed Share for switching between experts

Conceptual message:

- Adapting to changing environment is not automatic
- Modelling with explicit sequences