

Machine Learning Theory 2025

Lecture 10

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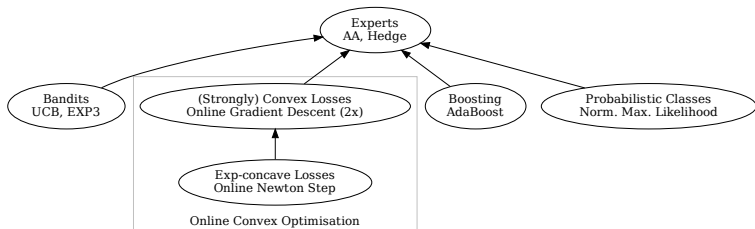
Online Convex Optimisation

- ▶ Gradient Descent for Convex Losses
- ▶ Online to Batch Conversion
- ▶ Gradient Descent for Strongly Convex Losses



Recap

Overview of Second Half of Course



Material: course notes on MLT website.

Recap: Finite Classes

So far we have seen learning “finite sets”:

Our learning algorithms behave like the **best** among K strategies.

- ▶ K -Experts setting
 - ▶ Mix loss : Aggregating Algorithm
 - ▶ Dot loss : Hedge algorithm
- ▶ K -armed bandit settings
 - ▶ Adversarial bandit : EXP3
 - ▶ Stochastic bandit : UCB

Outlook: Beyond the Finite

What if we want to compete with **infinite** sets?

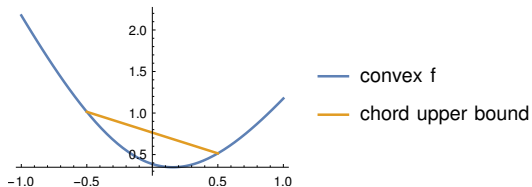
- ▶ Can we?
- ▶ How?

In each case, **lower bounds** grow with K : $\ln K$, $\sqrt{T \ln K}$, $\sqrt{TK \ln K}$, $K/\Delta \ln T$. So hopeless in the **unstructured** $K \rightarrow \infty$ case.

Today: compete with **continuous** sets of actions, parameterised such that the loss is a **convex** function of the action.

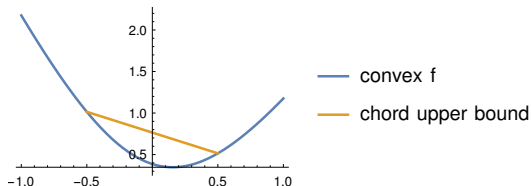
Convexity Review

Convex Functions I : definition



Fix a convex set $\mathcal{U} \subseteq \mathbb{R}^d$.

Convex Functions I : definition



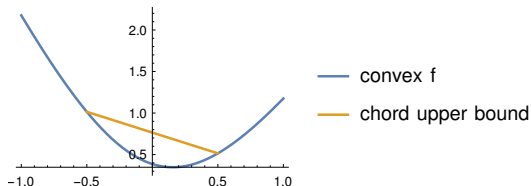
Fix a convex set $\mathcal{U} \subseteq \mathbb{R}^d$.

Definition

A function $f : \mathcal{U} \rightarrow \mathbb{R}$ is convex if for all $\mathbf{x}, \mathbf{y} \in \mathcal{U}$ and weights $\theta \in [0, 1]$,

$$f(\theta \mathbf{x} + (1 - \theta) \mathbf{y}) \leq \theta f(\mathbf{x}) + (1 - \theta) f(\mathbf{y}).$$

Convex Functions I : definition



Fix a convex set $\mathcal{U} \subseteq \mathbb{R}^d$.

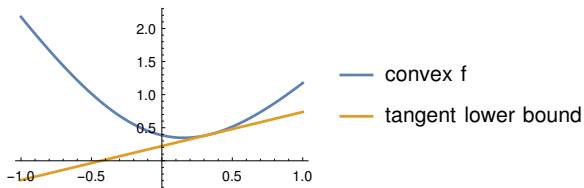
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Extends to arbitrary mixtures: $f(\mathbb{E}[X]) \leq \mathbb{E}[f(X)]$ (Jensen).

Convex Functions II : tangent bound

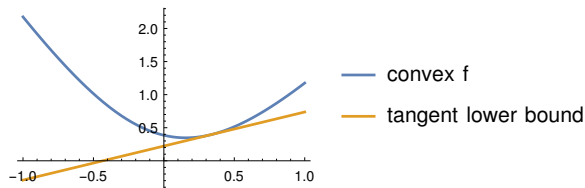


Fact

A differentiable function $f : \mathcal{U} \rightarrow \mathbb{R}$ is convex iff for all $x, y \in \mathcal{U}$

$$f(y) - f(x) \geq \langle y - x, \nabla f(x) \rangle$$

Convex Functions II : tangent bound



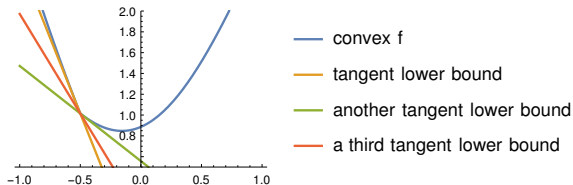
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Symmetrically, $\langle y - x, \nabla f(y) \rangle \geq f(y) - f(x)$.

Convex Functions III : sub-gradient



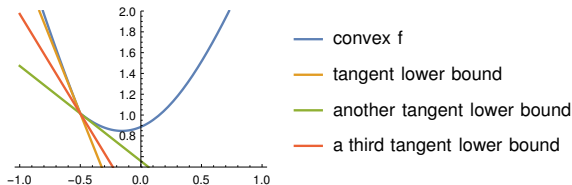
Fact (Sub-gradient)

For any convex $f : \mathcal{U} \rightarrow \mathbb{R}$, possibly non-differentiable, and point $x \in \mathcal{U}$, there always exists **some** vector $g \in \mathbb{R}^d$ such that for all $y \in \mathcal{U}$

$$f(y) - f(x) \geq \langle y - x, g \rangle$$

Any such vector g is called a **sub-gradient** (of f at x).

Convex Functions III : sub-gradient



Fact (Sub-gradient)

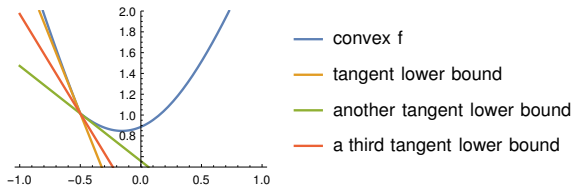
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The gradient of a differentiable function is a sub-gradient.

Convex Functions III : sub-gradient



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The gradient of a differentiable function is a sub-gradient.

We will abuse notation and denote **any** sub-gradient by $\nabla f(x)$.

Online Convex Optimisation

Online Convex Optimisation

General yet simple sequential decision problem.

Fix a convex set $\mathcal{U} \subseteq \mathbb{R}^d$.

Protocol

For $t = 1, 2, \dots$

- ▶ Learner chooses a point $\mathbf{w}_t \in \mathcal{U}$.
- ▶ Adversary reveals convex loss function $f_t : \mathcal{U} \rightarrow \mathbb{R}$.
- ▶ Learner's loss is $f_t(\mathbf{w}_t)$

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Objective:

Regret w.r.t. best point after T rounds:

$$R_T = \max_{\mathbf{u} \in \mathcal{U}} \sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{u}))$$

Example loss functions

Setting	loss function $f_t(\mathbf{u})$
Hedge setting	$\mathbf{u}^\top \ell_t$
Point prediction	$\ \mathbf{u} - \mathbf{x}_t\ ^2$
Regression	$(\mathbf{u}^\top \mathbf{x}_t - y_t)^2$
Logistic regression	$\ln(1 + e^{-y_t \mathbf{u}^\top \mathbf{x}_t})$
Hinge loss	$\max\{0, 1 - y_t \mathbf{u}^\top \mathbf{x}_t\}$
Investment	$-\ln(\mathbf{u}^\top \mathbf{x}_t)$
Offline optimisation	$f(\mathbf{u})$

Online Gradient Descent (OGD)

Let $\mathcal{U} \subseteq \mathbb{R}^d$ be a closed convex set containing $\mathbf{0}$.

Definition

Online Gradient Descent with learning rate $\eta > 0$ plays

$$\mathbf{w}_1 = \mathbf{0} \quad \text{and} \quad \mathbf{w}_{t+1} = \Pi_{\mathcal{U}}(\mathbf{w}_t - \eta \nabla f_t(\mathbf{w}_t))$$

where

$$\Pi_{\mathcal{U}}(\mathbf{w}) = \arg \min_{\mathbf{u} \in \mathcal{U}} \|\mathbf{u} - \mathbf{w}\|$$

is the projection of $\mathbf{w} \in \mathbb{R}^d$ onto \mathcal{U} in Euclidean norm.

Online Gradient Descent (OGD)

Assumption: Bounded Gradients and Domain

Let G and D bound the gradients and the domain, i.e.

$$\|\nabla f_t(\mathbf{u})\| \leq G \quad \text{and} \quad \|\mathbf{u}\| \leq D \quad \text{for all } \mathbf{u} \in \mathcal{U}.$$

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Theorem (OGD regret bound)

Online Gradient Descent guarantees

$$R_T = \max_{\mathbf{u} \in \mathcal{U}} \sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{u})) \leq \frac{1}{2\eta} D^2 + \frac{\eta}{2} T G^2$$

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Corollary

Tuning $\eta = \frac{D}{G\sqrt{T}}$ results in

$$R_T \leq D G \sqrt{T}$$

Pythagorean Inequality

Lemma (Pythagorean Inequality)

Fix a closed convex set $\mathcal{U} \subseteq \mathbb{R}^d$. Let $\mathbf{x} \in \mathcal{U}$, $\mathbf{y} \in \mathbb{R}^d$ and

$$\hat{\mathbf{y}} = \Pi_{\mathcal{U}}(\mathbf{y}) = \arg \min_{\mathbf{u} \in \mathcal{U}} \|\mathbf{u} - \mathbf{y}\|^2.$$

Then

$$\|\mathbf{x} - \hat{\mathbf{y}}\|^2 + \|\hat{\mathbf{y}} - \mathbf{y}\|^2 \leq \|\mathbf{x} - \mathbf{y}\|^2$$

NB: not to be confused with the **triangle inequality**

$$\|\mathbf{x} - \mathbf{y}\| \leq \|\mathbf{x} - \hat{\mathbf{y}}\| + \|\hat{\mathbf{y}} - \mathbf{y}\|.$$

Proof of GD regret bound I

Fix any $\mathbf{u} \in \mathcal{U}$. By convexity,

$$f_t(\mathbf{w}_t) - f_t(\mathbf{u}) \leq \langle \mathbf{w}_t - \mathbf{u}, \nabla f_t(\mathbf{w}_t) \rangle.$$

By the OGD update rule

$$\begin{aligned} \|\mathbf{w}_{t+1} - \mathbf{u}\|^2 &= \|\Pi_{\mathcal{U}}(\mathbf{w}_t - \eta \nabla f_t(\mathbf{w}_t)) - \mathbf{u}\|^2 \\ &\stackrel{\text{Pyth. Ineq.}}{\leq} \|\mathbf{w}_t - \eta \nabla f_t(\mathbf{w}_t) - \mathbf{u}\|^2 \\ &= \|\mathbf{w}_t - \mathbf{u}\|^2 - 2\eta \langle \mathbf{w}_t - \mathbf{u}, \nabla f_t(\mathbf{w}_t) \rangle + \eta^2 \|\nabla f_t(\mathbf{w}_t)\|^2. \end{aligned}$$

We can chain these two inequalities into

$$\begin{aligned} f_t(\mathbf{w}_t) - f_t(\mathbf{u}) &\leq \langle \mathbf{w}_t - \mathbf{u}, \nabla f_t(\mathbf{w}_t) \rangle \\ &\leq \frac{\|\mathbf{w}_t - \mathbf{u}\|^2 - \|\mathbf{w}_{t+1} - \mathbf{u}\|^2}{2\eta} + \frac{\eta}{2} \|\nabla f_t(\mathbf{w}_t)\|^2. \end{aligned}$$

Proof of GD regret bound II

Summing over T rounds, we find

$$\begin{aligned}\sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{u})) &\leq \underbrace{\sum_{t=1}^T \frac{\|\mathbf{w}_t - \mathbf{u}\|^2 - \|\mathbf{w}_{t+1} - \mathbf{u}\|^2}{2\eta}}_{\text{telescopes}} + \frac{\eta}{2} \sum_{t=1}^T \|\nabla f_t(\mathbf{w}_t)\|^2 \\ &\leq \frac{\|\mathbf{u}\|^2 - \|\mathbf{w}_{T+1} - \mathbf{u}\|^2}{2\eta} + \frac{\eta}{2} \sum_{t=1}^T \|\nabla f_t(\mathbf{w}_t)\|^2 \\ &\leq \frac{D^2}{2\eta} + \frac{\eta}{2} TG^2.\end{aligned}$$

Conclusion of Convex Losses part

We developed Online Gradient Descent.

OGD behaves almost as well as the **best point** in the continuous domain, as measured by a sum of **adversarially** chosen convex loss functions.

Observations:

- ▶ $GD\sqrt{T}$ regret
- ▶ Efficient: run time is $O(d)$ per round, plus
 - ▶ one gradient evaluation
 - ▶ one projection onto the domain

Online to Batch Conversion

Online to Batch Conversion

We revisit statistical learning, with loss functions abstracting examples:

Goal: obtain an estimator \hat{w}_T with small expected excess risk.

$$\mathbb{E}_{f_1, \dots, f_T} \left[\mathbb{E}_f [f(\hat{w}_T) - f(u^*)] \right] \leq \text{small}$$

where the **training set** f_1, \dots, f_T and the **test sample** f are drawn i.i.d. and u^* optimises the risk $u \mapsto \mathbb{E}_f[f(u)]$.

How to design \hat{w}_T ?

Online to Batch Conversion

Idea: use online learning algorithm. Given training sample f_1, \dots, f_T , the algorithm picks w_1, \dots, w_T . Let us define the *average iterate estimator*

$$\hat{w}_T = \frac{1}{T} \sum_{t=1}^T w_t.$$

Online to Batch Conversion

Idea: use online learning algorithm. Given training sample f_1, \dots, f_T , the algorithm picks w_1, \dots, w_T . Let us define the *average iterate estimator*

$$\hat{w}_T = \frac{1}{T} \sum_{t=1}^T w_t.$$

Theorem

An online regret bound $R_T \leq B(T)$ implies

$$\mathbb{E}_{iid\ f_1, \dots, f_T, f} [f(\hat{w}_T) - f(u^*)] \leq \frac{B(T)}{T}$$

Online to Batch Proof

$$\begin{aligned} & \mathbb{E}_{\text{iid } f_1, \dots, f_T, f} [f(\hat{w}_T) - f(u^*)] \\ & \leq \mathbb{E}_{\text{iid } f_1, \dots, f_T, f} \left[\frac{1}{T} \sum_{t=1}^T (f(w_t) - f(u^*)) \right] \\ & = \mathbb{E}_{\text{iid } f_1, \dots, f_T} \left[\frac{1}{T} \sum_{t=1}^T (f_t(w_t) - f_t(u^*)) \right] \leq \frac{B(T)}{T} \end{aligned}$$

The first step is convexity of f . The last step uses that f and f_t have the same distribution (and w_t is not a function of f_t).

Online to Batch Proof

$$\begin{aligned} & \mathbb{E}_{\text{iid } f_1, \dots, f_T, f} [f(\hat{w}_T) - f(u^*)] \\ & \leq \mathbb{E}_{\text{iid } f_1, \dots, f_T, f} \left[\frac{1}{T} \sum_{t=1}^T (f(w_t) - f(u^*)) \right] \\ & = \mathbb{E}_{\text{iid } f_1, \dots, f_T} \left[\frac{1}{T} \sum_{t=1}^T (f_t(w_t) - f_t(u^*)) \right] \leq \frac{B(T)}{T} \end{aligned}$$

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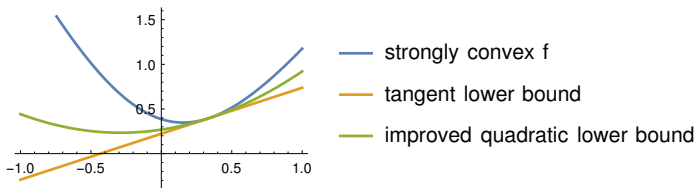
We can use online learning methods for statistical learning.

Online Strongly Convex Optimisation

Structure

What if we **know more** about my setting than **convexity of the loss function**? Can we learn faster?

Strongly Convex Case

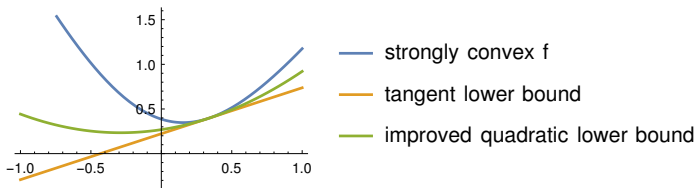


Definition

A function $f : \mathcal{U} \rightarrow \mathbb{R}$ is *strongly convex* to degree $\alpha \geq 0$ if

$$f(u) - f(w) \geq \langle u - w, \nabla f(w) \rangle + \frac{\alpha}{2} \|u - w\|^2$$

Strongly Convex Case



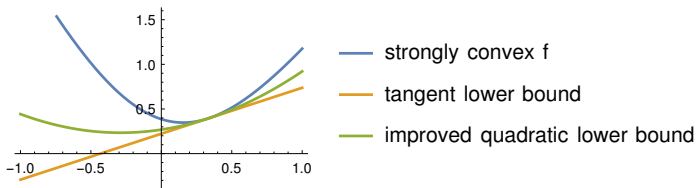
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Example: $f(w) = \frac{1}{2} \|w - x_t\|^2$ is strongly convex with $\alpha = 1$.

Strongly Convex Case



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Example: $f(w) = \frac{1}{2} \|w - x_t\|^2$ is strongly convex with $\alpha = 1$.

Idea: could this extra knowledge help in the regret rate?

Online Gradient Descent with time-varying learning rate

Definition (OGD with time-varying learning rate)

$$\boldsymbol{w}_1 = \mathbf{0} \quad \text{and} \quad \boldsymbol{w}_{t+1} = \Pi_{\mathcal{U}}(\boldsymbol{w}_t - \eta_t \nabla f_t(\boldsymbol{w}_t))$$

Online Gradient Descent with time-varying learning rate

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Theorem

For α -strongly convex loss functions, OGD with learning rate $\eta_t = \frac{1}{\alpha t}$ ensures

$$R_T \leq \frac{G^2}{2\alpha} (1 + \ln T).$$

Proof I

Exactly as for the convex case, the update rule ensures

$$\begin{aligned}\|\mathbf{w}_{t+1} - \mathbf{u}\|^2 &= \|\Pi_{\mathcal{U}}(\mathbf{w}_t - \eta_t \nabla f_t(\mathbf{w}_t)) - \mathbf{u}\|^2 \\ &\stackrel{\text{Pyth. Ineq.}}{\leq} \|\mathbf{w}_t - \eta_t \nabla f_t(\mathbf{w}_t) - \mathbf{u}\|^2 \\ &= \|\mathbf{w}_t - \mathbf{u}\|^2 - 2\eta_t \langle \mathbf{w}_t - \mathbf{u}, \nabla f_t(\mathbf{w}_t) \rangle + \eta_t^2 \|\nabla f_t(\mathbf{w}_t)\|^2\end{aligned}$$

Combination with strong convexity gives

$$\begin{aligned}&f_t(\mathbf{w}_t) - f_t(\mathbf{u}) \\ &\leq \langle \mathbf{w}_t - \mathbf{u}, \nabla f_t(\mathbf{w}_t) \rangle - \frac{\alpha}{2} \|\mathbf{w}_t - \mathbf{u}\|^2 \\ &\leq \frac{\|\mathbf{w}_t - \mathbf{u}\|^2 - \|\mathbf{w}_{t+1} - \mathbf{u}\|^2 + \eta_t^2 \|\nabla f_t(\mathbf{w}_t)\|^2}{2\eta_t} - \frac{\alpha}{2} \|\mathbf{w}_t - \mathbf{u}\|^2 \\ &= \|\mathbf{w}_t - \mathbf{u}\|^2 \left(\frac{1}{2\eta_t} - \frac{\alpha}{2} \right) - \frac{\|\mathbf{w}_{t+1} - \mathbf{u}\|^2}{2\eta_t} + \frac{\eta_t \|\nabla f_t(\mathbf{w}_t)\|^2}{2}\end{aligned}$$

Proof II

Summing over rounds gives

$$\begin{aligned} & \sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{u})) \\ & \leq \sum_{t=1}^T \left(\|\mathbf{w}_t - \mathbf{u}\|^2 \left(\frac{1}{2\eta_t} - \frac{\alpha}{2} \right) - \frac{\|\mathbf{w}_{t+1} - \mathbf{u}\|^2}{2\eta_t} + \frac{\eta_t \|\nabla f_t(\mathbf{w}_t)\|^2}{2} \right) \\ & = \|\mathbf{w}_1 - \mathbf{u}\|^2 \left(\frac{1}{2\eta_1} - \frac{\alpha}{2} \right) + \sum_{t=2}^T \|\mathbf{w}_t - \mathbf{u}\|^2 \left(\frac{1}{2\eta_t} - \frac{\alpha}{2} - \frac{1}{2\eta_{t-1}} \right) \\ & \quad - \frac{\|\mathbf{w}_{T+1} - \mathbf{u}\|^2}{2\eta_T} + \sum_{t=1}^T \frac{\eta_t \|\nabla f_t(\mathbf{w}_t)\|^2}{2} \end{aligned}$$

Key idea for telescoping is to cancel coefficient on $\|\mathbf{w}_t - \mathbf{u}\|^2$ in the sum:

$$\frac{1}{2\eta_t} - \frac{\alpha}{2} - \frac{1}{2\eta_{t-1}} = 0.$$

Proof III

This yields recurrence

$$\eta_t = \frac{1}{\frac{1}{\eta_{t-1}} + \alpha}$$

Cancelling the coefficient on $\|\mathbf{w}_1 - \mathbf{u}\|^2$ gives starting point $\eta_1 = \frac{1}{\alpha}$.
This leads to overall solution $\eta_t = \frac{1}{\alpha t}$. Plugging that in, we find

$$\sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{u})) \leq \sum_{t=1}^T \frac{\|\nabla f_t(\mathbf{w}_t)\|^2}{2\alpha t} \leq \frac{G^2}{2\alpha} (1 + \ln T).$$

Conclusion

Tools for learning in convex settings.

- ▶ Guaranteed robustness against adversarial losses
- ▶ Efficient
- ▶ Building block for
 - ▶ Learning in non-convex settings (AdaGrad for DNN)
 - ▶ Learning in games
 - ▶ Non-convex games (GANs)
 - ▶ ...