

Quasi-Measurable Spaces

A Convenient Foundation of Probability Theory

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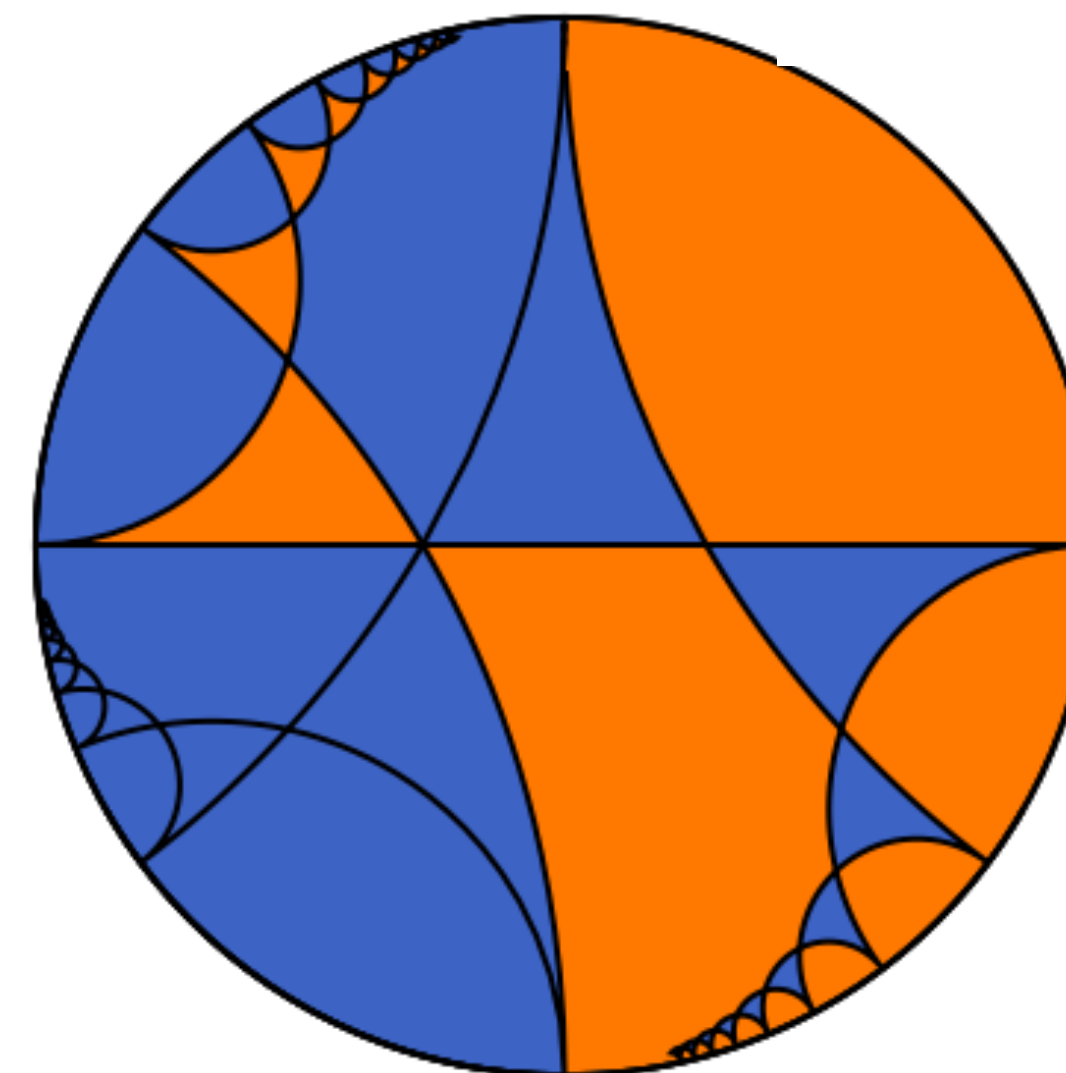
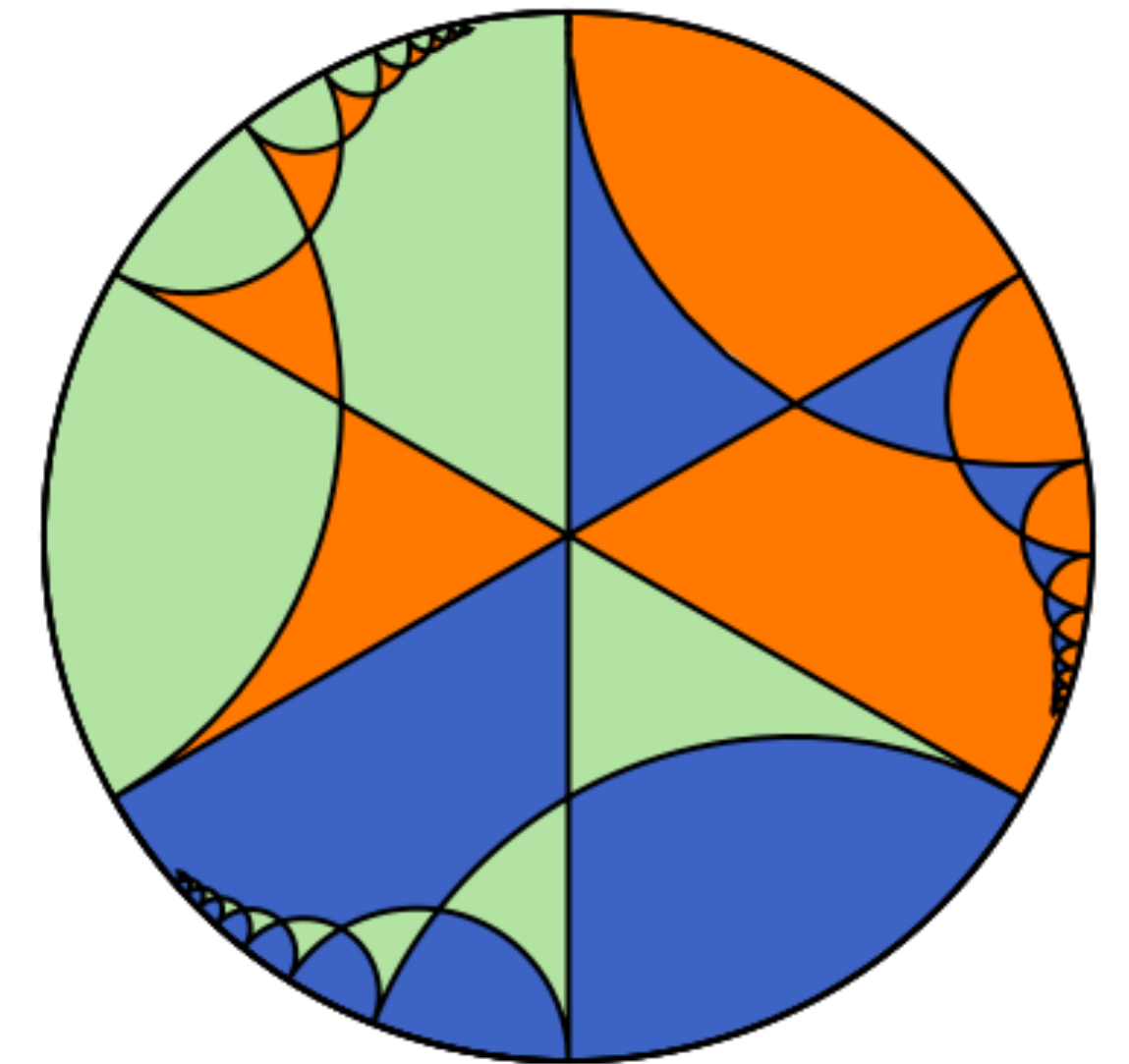
References

- The talk is based on the following papers:
 - Chris Heunen, Ohad Kammar, Sam Staton, Hongseok Yang. ***A convenient category for higher-order probability theory.*** 32nd Annual ACM/IEEE Symposium on Logic in Computer Science (LICS), 2017.
 - Patrick Forré, ***Quasi-Measurable Spaces***, 2021, <https://arxiv.org/abs/2109.11631>.

Why Measure Theory to do Probability Theory?

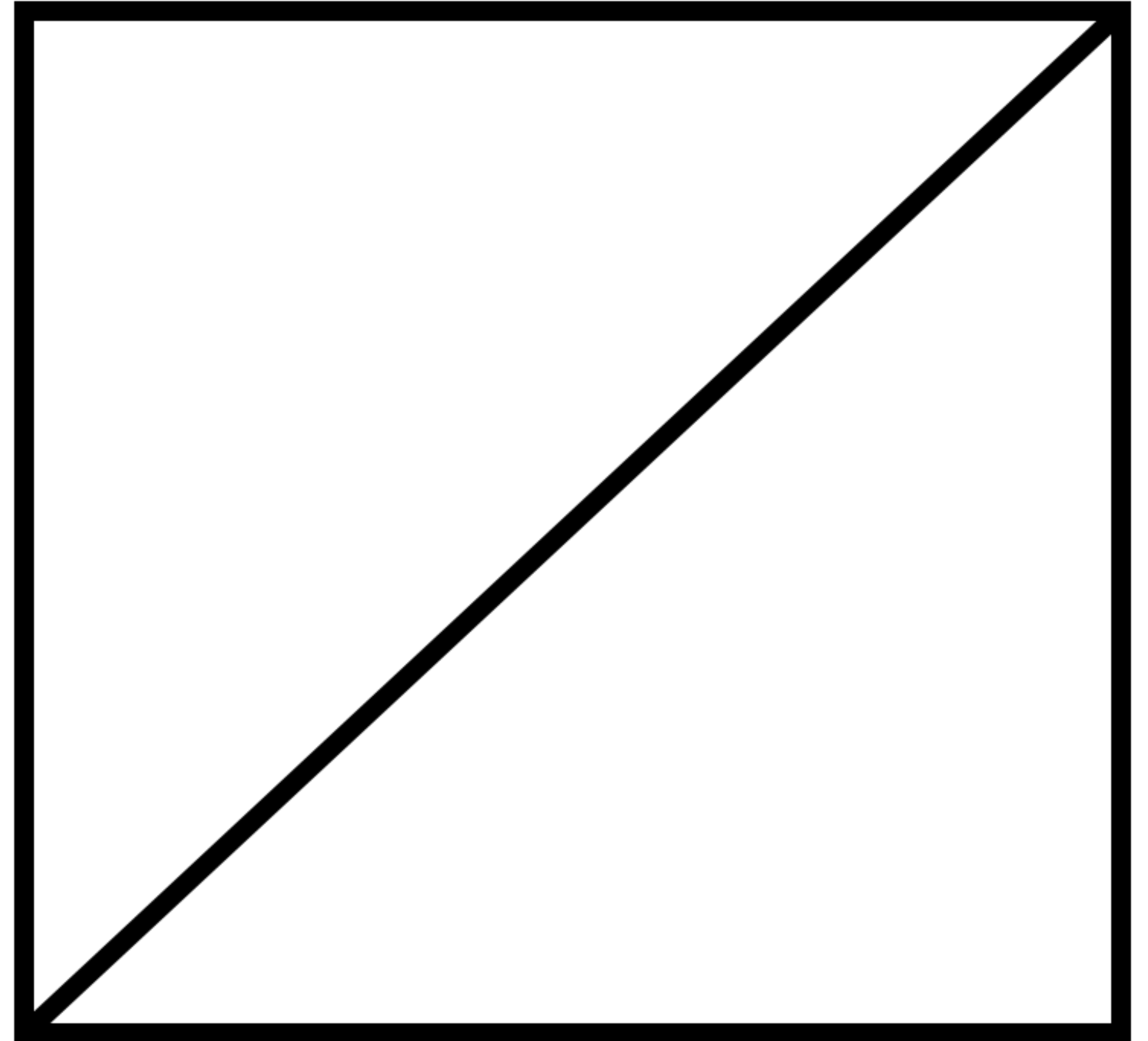
Why Measure Theory in the first place?

- The existence of the Lebesgue measure:
 - does not exist on whole power set $2^{\mathbb{R}}$.
 - but does exist on Borel σ -algebra $\mathcal{B}_{\mathbb{R}}$.
- To prevent set-theoretic paradoxa like Banach-Tarski:
 - the orange is both a third and a half of the Poincaré disk / hyperbolic plane.



Discrete and continuous distributions are not expressive enough

- The uniform distribution on the diagonal $\Delta \subseteq [0,1]^2$
 - is neither discrete nor absolute continuous w.r.t. λ^2 .
 - so it can not be described with a probability mass function nor with a probability density w.r.t. λ^2 .



The Category of Measurable Spaces

- Let \mathcal{X} be a set. A σ -algebra on \mathcal{X} is a set of subsets $\mathcal{B} \subseteq 2^{\mathcal{X}}$ such that:
 - $\emptyset \in \mathcal{B}$,
 - $A_n \in \mathcal{B}, n \in \mathbb{N} \implies \bigcup_{n \in \mathbb{N}} A_n \in \mathcal{B}$
 - $A \in \mathcal{B} \implies \mathcal{X} \setminus A \in \mathcal{B}$
- A tuple $(\mathcal{X}, \mathcal{B}_{\mathcal{X}})$ of a set \mathcal{X} and a σ -algebra $\mathcal{B}_{\mathcal{X}}$ is called **measurable space**.
- A map $f: \mathcal{X} \rightarrow \mathcal{Y}$ between measurable spaces $(\mathcal{X}, \mathcal{B}_{\mathcal{X}})$ and $(\mathcal{Y}, \mathcal{B}_{\mathcal{Y}})$ is called a **measurable map** if:
$$B \in \mathcal{B}_{\mathcal{Y}} \implies f^{-1}(B) \in \mathcal{B}_{\mathcal{X}}.$$
 - Note that the compositions of two measurable maps is a measurable map.
- Meas denotes the **category of measurable spaces and measurable maps**.

Kolmogorov's approach to Probability Theory (1933)

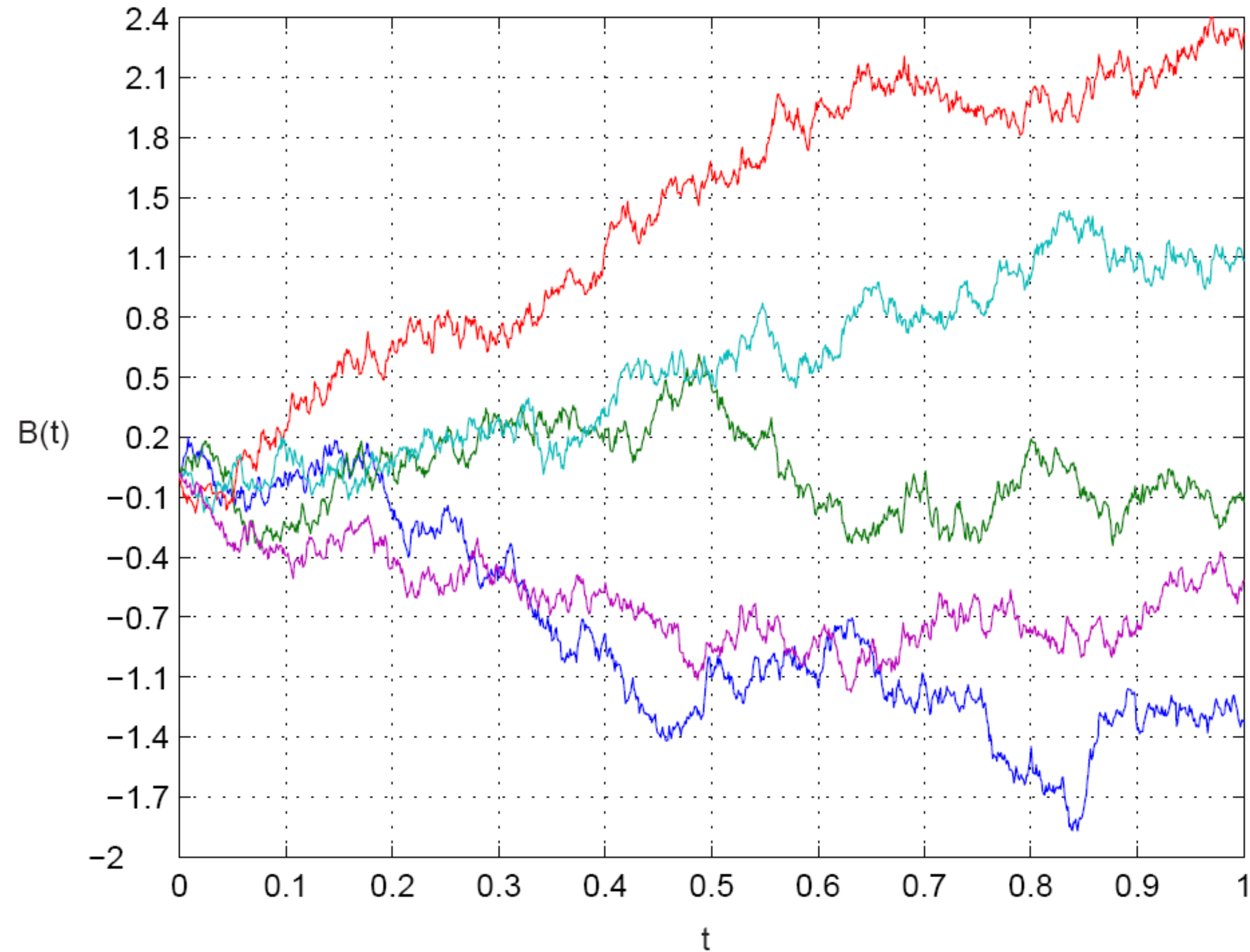
- **Kolmogorov Axioms:**
 - A probability distribution is just a normalized measure.
- Probability Theory can thus be viewed as a sub-field of Measure Theory.
 - now allows for Lebesgue's theory of integration (measure integrals, etc.)
- Measure Theory as an expressive "safe space" of Probability Theory.
- Andrei Kolmogoroff. *Grundbegriffe der Wahrscheinlichkeitsrechnung*. Ergebnisse der Mathematik und Ihrer Grenzgebiete. 1. Folge, Nr. 2, Springer (1933).

How to formalize Random Variables?

- **Sample space** is a measurable space: $(\Omega, \mathcal{B}_\Omega)$
 - where \mathcal{B}_Ω is the σ -algebra/set of admissible outcome events on Ω .
- **State space** is a measurable space: $(\mathcal{X}, \mathcal{B}_\mathcal{X})$
 - where $\mathcal{B}_\mathcal{X}$ is another σ -algebra/set of admissible events on \mathcal{X} .
- Admissible **random variables** are all measurable maps:
 - $X \in \text{Meas}((\Omega, \mathcal{B}_\Omega), (\mathcal{X}, \mathcal{B}_\mathcal{X}))$
- For fixed **probability measure** P on $(\Omega, \mathcal{B}_\Omega)$ the distribution of X is:
 - push-forward probability measure: X_*P on $\mathcal{B}_\mathcal{X}$ (also written as: $P(X)$).

What is a Stochastic Process?

Different realizations of one stochastic process



Text book definitions - stochastic process

- *D. Revus, M. Yor - Continuous Martingales and Brownian Motion:*

(1.1) Definition. *Let T be a set, (E, \mathcal{E}) a measurable space. A stochastic process indexed by T , taking its values in (E, \mathcal{E}) , is a family of measurable mappings $X_t, t \in T$, from a probability space (Ω, \mathcal{F}, P) into (E, \mathcal{E}) . The space (E, \mathcal{E}) is called the state space.*

For every $\omega \in \Omega$, the mapping $t \rightarrow X_t(\omega)$ is a “curve” in E which is referred to as a *trajectory* or a *path* of X . We may think of a path as a point chosen randomly in the space $\mathcal{F}(T, E)$ of all functions from T into E , or, as we shall see later, in a reasonable subset of this space.

Text book definitions - stochastic process

- *D. Revus, M. Yor - Continuous Martingales and Brownian Motion:*

Let (\mathcal{F}_t) be a filtration on (Ω, \mathcal{F}) and T a stopping time. For a process X , we define a new mapping X_T on the set $\{\omega : T(\omega) < \infty\}$ by

$$X_T(\omega) = X_t(\omega) \quad \text{if} \quad T(\omega) = t.$$

This is the position of the process X at time T , but it is not clear that X_T is a random variable on $\{T < \infty\}$. Moreover if X is adapted, we would like X_T

(4.7) Definition. *A process X is progressively measurable or simply progressive (with respect to the filtration (\mathcal{F}_t)) if for every t the map $(s, \omega) \rightarrow X_s(\omega)$ from $[0, t] \times \Omega$ into (E, \mathcal{E}) is $\mathcal{B}([0, t]) \otimes \mathcal{F}_t$ -measurable. A subset Γ of $\mathbb{R}_+ \times \Omega$ is progressive if the process $X = 1_\Gamma$ is progressive.*

(4.8) Proposition. *An adapted process with right or left continuous paths is progressively measurable.*

Three Definitions of Stochastic Processes

- Let $(\Omega, \mathcal{B}_\Omega)$ be the sample space, $(\mathcal{X}, \mathcal{B}_\mathcal{X})$ the state space, $(\mathcal{T}, \mathcal{B}_\mathcal{T})$ the time space, e.g. $\mathcal{T} = \mathbb{N}$ or $\mathcal{T} = \mathbb{R}_{\geq 0}$.
- A *stochastic process* is what kind of random variable / measurable map?

$$1. \quad (X_t)_{t \in \mathcal{T}} : \Omega \rightarrow \prod_{t \in \mathcal{T}} \mathcal{X}, \quad \omega \mapsto (X_t(\omega))_{t \in \mathcal{T}},$$

$$2. \quad X : \Omega \rightarrow \text{Meas}(\mathcal{T}, \mathcal{X}), \quad \omega \mapsto (t \mapsto X(\omega)(t)),$$

$$3. \quad X : \Omega \times \mathcal{T} \rightarrow \mathcal{X}, \quad (\omega, t) \mapsto X(\omega, t).$$

Problems

- Three different (partially inconsistent) definitions of stochastic processes
 - mismatch between formalization and meaning.
- Not clear how to turn $\text{Meas}(\mathcal{T}, \mathcal{X})$ into a measurable space in itself?
 - existence of well-behaved σ -algebra unclear
- The measurability of $\omega \mapsto X_{T(\omega)}(\omega)$ only guaranteed under additional assumptions.

Are Probabilistic Programs functional?

A program that outputs a probabilistic program

```
def prog_prob_prog(a):  
    return lambda m,s: [Z:=np.random.uniform(), a*m+s*Z][-1]
```

```
print(prog_prob_prog(a=1))
```

```
<function prog_prob_prog.<locals>.<lambda>
```

```
for n in range(5):  
    print(prog_prob_prog(a=n)(m=5,s=2))
```

```
1.8106662116099772  
6.762509413168864  
10.365457994333775  
15.884402920590935  
21.48676872656254
```

- uncurried version:

```
def prob_prog(a,m,s):  
    Z=np.random.uniform()  
    return a*m+s*Z
```

```
for n in range(5):  
    print(prob_prog(a=n,m=5,s=2))
```

```
1.5413066059310134  
6.248544376268809  
10.923467140491365  
16.8296978388216  
20.72122425243884
```


Another probabilistic program that outputs a probabilistic program

```
import numpy as np

def prob_prog_1(m=0, s=1):
    Z = np.random.normal()
    X = m+s*Z
    return(X)
```

- Output:

```
for n in range(10):
    print(prob_prog_2(7,3,2))
```

```
('output', 13.334218167868446)
('output', 11.471183953689039)
('det fct', <function prob_prog_2.<locals>.<lambda> at 0x7...
('prob_prog_1', <function prob_prog_1 at 0x7fd4f2820550>)
('prob_prog_1', <function prob_prog_1 at 0x7fd4f2820550>)
('prob_prog_1', <function prob_prog_1 at 0x7fd4f2820550>)
('prob_prog_1', <function prob_prog_1 at 0x7fd4f2820550>)
('output', 8.377835031430754)
('prob_prog_1', <function prob_prog_1 at 0x7fd4f2820550>)
('output', 11.669097814447499)
```

```
def prob_prog_2(m=0, s=1, b=0):
    U = np.random.uniform()
    if U <= 0.33:
        return 'det fct', lambda x: x**2
    elif U <= 0.66:
        return 'output', b+prob_prog_1(m, s)
    else:
        return 'prob_prog_1', prob_prog_1
```

- How can one mathematically describe such a (probabilistic) program that outputs a probabilistic program, possibly of different types?

How to formalize Probabilistic Programs?

- Probabilistic programs:
 - take input x ,
 - sample internal random number ω ,
 - determine (stochastic) output z ,
- so either:
 - measurable map: $K : \Omega \times \mathcal{X} \rightarrow \mathcal{Z}$,
 - however, $\omega \in \Omega$ not really an input, rather internal
 - Markov kernels from input space \mathcal{X} to output space \mathcal{Z} ,
 - measurable map $K : \mathcal{X} \rightarrow \mathcal{P}(\mathcal{Z})$
 - set of probabilistic programs: $\text{Meas}(\mathcal{X}, \mathcal{P}(\mathcal{Z}))$.

Church's Simply Typed λ -Calculus (1940)

- **Functional Programming** should satisfy “**curry**” / “**uncurry**” operations:
 - $(x, y) \mapsto f(x, y)$ corresponds 1:1 to: $x \mapsto (y \mapsto f(x, y))$
 - $(x, y) \mapsto g(x)(y)$ corresponds 1:1 to: $x \mapsto g(x) = (y \mapsto g(x)(y))$
- This mean a program in two (or more) variables $f(x, y)$ can be expressed as iteratively defining a functions in one variable $g(x)(y)$ and vice versa.
- Requires program-valued programs / function-valued functions.
- Realized in functional programming language Haskell.
- Mathematically corresponds to cartesian closed categories.
- Alonzo Church. *A formulation of the simple theory of types*. The Journal of Symbolic Logic 5.2 (1940): 56-68.

Can we Curry / Uncurry Probabilistic Programs?

- Curry / Uncurry operations would translate to isomorphism:
 - $\text{Meas} \left(\mathcal{X} \times \mathcal{Y}, \mathcal{P}(\mathcal{Z}) \right) \cong \text{Meas} \left(\mathcal{X}, \text{Meas} \left(\mathcal{Y}, \mathcal{P}(\mathcal{Z}) \right) \right)$.
- This means we need to be able to mathematically describe programs whose outputs are probabilistic programs.
- Furthermore, we need the operation \mathcal{P} to be well-behaved:
 - functorial, respects product structure
 - strong probability monad

Remark - Monad

- **Monad**
 - theory of functional programming with side effects
 - equivalent to category-theoretical construction in mathematics
- **strong monad:**
 - well-behaved w.r.t. products

Definition - Monad

- A **Monad** on a category \mathcal{C} is a triple $(\mathcal{P}, \delta, \mathbb{M})$ consisting of:
 - a functor $\mathcal{P} : \mathcal{C} \rightarrow \mathcal{C}$,
 - a natural transformation $\delta : \text{id}_{\mathcal{C}} \rightarrow \mathcal{P}$,
 - a natural transformation $\mathbb{M} : \mathcal{P}^2 := \mathcal{P} \circ \mathcal{P} \rightarrow \mathcal{P}$,
- such that:
 - $\mathbb{M} \circ \mathcal{P}\mathbb{M} = \mathbb{M} \circ \mathbb{M}\mathcal{P}$ as natural transformations $\mathcal{P}^3 \rightarrow \mathcal{P}$,
 - $\mathbb{M} \circ \mathcal{P}\delta = \mathbb{M} \circ \delta\mathcal{P} = \text{id}_{\mathcal{P}}$ as natural transformations $\mathcal{P} \rightarrow \mathcal{P}$.
- A monad is called **strong**, if it is also “well-behaved” w.r.t. finite products \times .

Problems

- Define strong probability monad $(\mathcal{P}, \delta, \mathbb{M})$
 - Giry monad defined on category of measurable spaces \mathbf{Meas} .
- The set of all programs that output probabilistic programs should be:
 - $\mathbf{Meas} \left(\mathcal{X}, \mathbf{Meas} \left(\mathcal{Y}, \mathcal{P}(\mathcal{I}) \right) \right)$
- Not clear how to turn $\mathbf{Meas} \left(\mathcal{Y}, \mathcal{P}(\mathcal{I}) \right)$ into a measurable space in itself
 - existence of well-behaved σ -algebra unclear
- Is it possible to do more complicated constructions, e.g. dependent products, etc,?
 - Can we get a dependent type theory together with higher-order probability theory?

**Can we do Graphical Reasoning between
Random Variables and Mechanisms?**

Conditional Independence in Probabilistic Graphical Models

- Consider a **Markov chain**:



- We have:

- factorization: $P(X, Y, Z) = P(Z|Y) \otimes P(Y|X) \otimes P(X)$

- tells us that Z is only dependent on Y , and, independent of X when conditioned on Y , but then also of the choice of $P(Y|X)$ and $P(X)$.

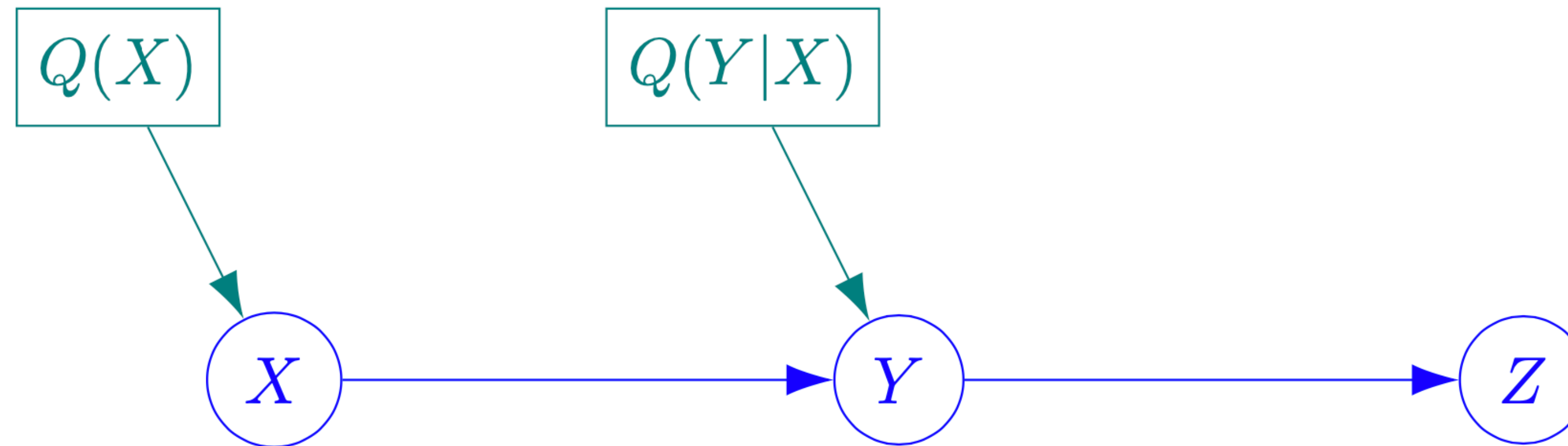
- We want to be able to:

- formalize conditional independence: $Z \perp\!\!\!\perp X, Q(Y|X), Q(X) | Y$

- including non-random variables $Q(X)$ and $Q(Y|X)$

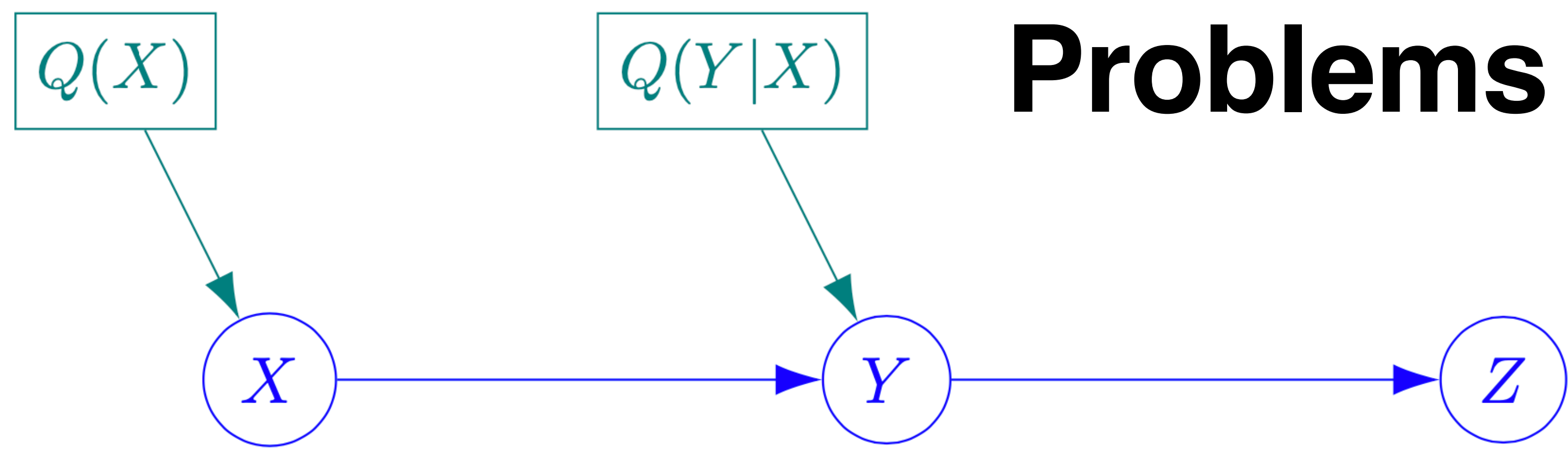
- read this off a graph via d-separation (or similar).

Including Non-Random Variables



- $Q(Y|X)$ is non-random and takes values in $\mathcal{L} := \text{Meas}(\mathcal{X}, \mathcal{P}(\mathcal{Y}))$
- Then Y is determined by the new mechanism:
 - $\mathcal{L} \times \mathcal{X} \rightarrow \mathcal{P}(\mathcal{Y}), \quad (Q(Y|X), x) \mapsto Q(Y|X = x).$
- similarly for X .

Problems



- Problems:
 - Not clear how to deal with non-random variables.
 - Not clear how to turn $\mathcal{L} := \text{Meas}(\mathcal{X}, \mathcal{P}(\mathcal{Y}))$ into a measurable space
 - Not clear how to define conditional independence with the two problems above.
 - Not clear if this corresponds to graphical conditional independence criteria.

How to formalize standard Causal Assumptions?

Causal Inference - Estimating Treatment Effects

- For estimating treatment effect, in the typical case, we have the variables:
 - X = observed treatment variable,
 - Y = observed outcome,
 - Y_x = potential outcome variable under (forced) treatment $X = x$,
 - Z = all other relevant features of the patient.
- Estimation is not possible without further assumptions.
- Typical assumptions made are:
 - Strong Ignorability: $X \perp\!\!\!\perp (Y_x)_{x \in \mathcal{X}} \mid Z$,
 - Consistency: $Y = Y_X$ a.s.

Problems

- Here, $(Y_x)_{x \in \mathcal{X}}$ is used as a vector of random variables from which we can pick components: $(\tilde{x}, (Y_x)_{x \in \mathcal{X}}) \mapsto Y_{\tilde{x}}$.

- However, the following map is, in general, not measurable:

$$\bullet \mathcal{X} \times \prod_{x \in \mathcal{X}} \mathcal{Y} \rightarrow \mathcal{Y}, \quad (\tilde{x}, (y_x)_{x \in \mathcal{X}}) \mapsto y_{\tilde{x}}.$$

- So Strong Ignorability does formally not go well with Consistency.
- Not immediate clear how to fix this.

Explanation

- Let $\mathcal{X} = \mathbb{R}$, $\mathcal{Y}_x := \mathcal{Y} := \{0,1\}$, then the following map is **not measurable**:

$$e : \mathcal{X} \times \prod_{x \in \mathcal{X}} \mathcal{Y} \rightarrow \mathcal{Y}, \quad (\tilde{x}, (y_x)_{x \in \mathcal{X}}) \mapsto y_{\tilde{x}}.$$

- Otherwise, we had: $D := e^{-1}(1) \in \mathcal{B}_{\mathcal{X}} \otimes \bigotimes_{x \in \mathcal{X}} \mathcal{B}_{\mathcal{Y}_x}$.

- Then D lies in a sub- σ -algebra generated by only countably many cylinder sets.

- So there exists countable subset: $\mathcal{C} \subseteq \mathcal{X}$ s.t.: $D = B \times \prod_{x \in \mathcal{X} \setminus \mathcal{C}} \mathcal{Y}_x$ with $B \subseteq \mathcal{X} \times \prod_{x \in \mathcal{C}} \mathcal{Y}_x$.

- For $x \in \mathcal{X} \setminus \mathcal{C}$: $e(x, 0_{\mathcal{C}}, 0_{\mathcal{X} \setminus \mathcal{C}}, 1_x, 0_{\mathcal{X} \setminus \mathcal{C}}) = 1$, so $(x, 0_{\mathcal{C}}, 0_{\mathcal{X} \setminus \mathcal{C}}, 1_x, 0_{\mathcal{X} \setminus \mathcal{C}}) \in D$, so $(x, 0_{\mathcal{C}}) \in B$.

- But then: $(x, 0_{\mathcal{C}}, 0_{\mathcal{X} \setminus \mathcal{C}}, 0_x, 0_{\mathcal{X} \setminus \mathcal{C}}) \in D = e^{-1}(1)$ and thus: $e(x, 0_{\mathcal{C}}, 0_{\mathcal{X} \setminus \mathcal{C}}, 0_x, 0_{\mathcal{X} \setminus \mathcal{C}}) = 1$.

- but clearly: $e(x, 0_{\mathcal{C}}, 0_{\mathcal{X} \setminus \mathcal{C}}, 0_x, 0_{\mathcal{X} \setminus \mathcal{C}}) = 0$, which is a contradiction.

How to formalize Counterfactual Probabilities?

Counterfactual Probabilities

- For reasoning about treatment effect we consider the variables:
 - X = observed treatment variable,
 - Y = observed outcome,
 - Y_x = potential outcome variable under (forced) treatment $X = x$.
- **Conditional counterfactual probabilities:**
 - $C(A | x, x') := P(Y_x \in A | X = x')$
 - “What would have happened (with which probability) under treatment $X = x$ given that the patient was actually treated with $X = x'$?”

Problems

- Not clear if conditional counterfactual probabilities are probability measures in A and/or measurable in x, x' or jointly.
 - $C(A | x, x') := P(Y_x \in A | X = x')$
- Not clear if conditioning is well-defined here, dependent on how to view $x \mapsto Y_x$.

**If we are going to change all of this
are we still able to do standard thing in
Probability Theory and (Bayesian) Statistics?**

Bad News

Random Functions do not exist in Meas

- Theorem (Aumann, 1961):
 - There is **no** σ -algebra $\mathcal{B}_{\mathcal{L}}$ on $\mathcal{L} := \text{Meas}(\mathbb{R}, \mathbb{R})$ such that the evaluation map is measurable:
 - $ev : \mathcal{L} \times \mathbb{R} \rightarrow \mathbb{R}, (f, x) \mapsto f(x),$
 - where \mathbb{R} carries the Borel- σ -algebra and \mathcal{L} is the space of all measurable maps from \mathbb{R} to \mathbb{R} , and the product carries the product- σ -algebra.
 - So there is no well-behaved way to define a probability distribution over all measurable functions in a fully non-parametric way.
- Robert J. Aumann. *Borel structures for function spaces*. Illinois Journal of Mathematics 5.4 (1961): 614-630.

Quasi-Measurable Spaces

Recall: Usual measure-theoretic approach

- **Sample space** is a measurable space: $(\Omega, \mathcal{B}_\Omega)$
 - where \mathcal{B}_Ω is the σ -algebra/set of admissible outcome events on Ω .
- **State space** is a measurable space: $(\mathcal{X}, \mathcal{B}_\mathcal{X})$
 - where $\mathcal{B}_\mathcal{X}$ is another σ -algebra/set of admissible events on \mathcal{X} .
- Admissible **random variables** are all measurable maps:
 - $X \in \text{Meas}((\Omega, \mathcal{B}_\Omega), (\mathcal{X}, \mathcal{B}_\mathcal{X}))$
- For fixed **probability measure** P on $(\Omega, \mathcal{B}_\Omega)$ the distribution of X is:
 - push-forward probability measure: X_*P on $\mathcal{B}_\mathcal{X}$ (also written as: $P(X)$).

Main Idea behind Quasi-Measurable Spaces

- Main idea: *Exchange the role of σ -algebras and random variables!!!*
- **Sample space** is a measurable space: $(\Omega, \mathcal{B}_\Omega)$
 - where \mathcal{B}_Ω is the σ -algebra/set of admissible outcome events on Ω .
- **State space** is a “*quasi-measurable space*”: $(\mathcal{X}, \mathcal{X}^\Omega)$
 - where \mathcal{X}^Ω is a set of admissible **random variables**.
- **σ -algebra** of admissible events is:
 - $\mathcal{B}_\mathcal{X} := \mathcal{B}(\mathcal{X}^\Omega) := \{A \subseteq \mathcal{X} \mid \forall X \in \mathcal{X}^\Omega . X^{-1}(A) \in \mathcal{B}_\Omega\}$
- For fixed **probability measure** P on $(\Omega, \mathcal{B}_\Omega)$ the distribution of X is:
 - push-forward probability measure: X_*P on $\mathcal{B}_\mathcal{X}$ (also written as: $P(X)$).

The Sample Space - Act 1 - Random Variables

- The **Sample Space** (Ω, Ω^Ω) consists of:
 - a set: Ω
 - a set of maps: $\Omega^\Omega \subseteq \{\Phi : \Omega \rightarrow \Omega\}$
 - such that:
 - $\text{id}_\Omega \in \Omega^\Omega$,
 - Ω^Ω contains all constant maps,
 - Ω^Ω is closed under composition:
 - $\Phi_1, \Phi_2 \in \Omega^\Omega \implies \Phi_2 \circ \Phi_1 \in \Omega^\Omega$.
- Standard example:
 - $\Omega^\Omega := \text{Meas} \left((\Omega, \mathcal{B}_\Omega), (\Omega, \mathcal{B}_\Omega) \right)$ for some carefully chosen σ -algebra: \mathcal{B}_Ω .

Quasi-Measurable Spaces

- A **Quasi-Measurable Space** $(\mathcal{X}, \mathcal{X}^\Omega)$ w.r.t. sample space (Ω, Ω^Ω) - per definition - consists of:
 - a set: \mathcal{X}
 - a set of **admissible random variables**: \mathcal{X}^Ω ,
 - i.e. a set of maps: $X : \Omega \rightarrow \mathcal{X}$, such that:
 - all *constant maps* $\Omega \rightarrow \mathcal{X}$ are in \mathcal{X}^Ω ,
 - \mathcal{X}^Ω is *closed under pre-composition* with Ω^Ω :
 - $X \in \mathcal{X}^\Omega, \Phi \in \Omega^\Omega \implies X \circ \Phi \in \mathcal{X}^\Omega$.

Quasi-Measurable Maps

- Let $(\mathcal{F}, \mathcal{F}^\Omega)$ and $(\mathcal{X}, \mathcal{X}^\Omega)$ two quasi-measurable spaces.
- A map $g : \mathcal{F} \rightarrow \mathcal{X}$ is called **quasi-measurable** if
 - $Z \in \mathcal{F}^\Omega \implies g(Z) := g \circ Z \in \mathcal{X}^\Omega$
- The set of all quasi-measurable maps is abbreviated:
 - $\text{QMS} \left((\mathcal{F}, \mathcal{F}^\Omega), (\mathcal{X}, \mathcal{X}^\Omega) \right)$ or $\text{QMS} (\mathcal{F}, \mathcal{X})$ for short.
- Note that the *composition* of two quasi-measurable maps is again *quasi-measurable*.
- The class of all quasi-measurable spaces (w.r.t. a fixed sample space) together with all **quasi-measurable maps** builds a **category**: QMS.

The Product Space

- Let $(\mathcal{X}_i, \mathcal{X}_i^\Omega)$ be a family of quasi-measurable spaces, $i \in I$.

- Then we turn the product space: $\prod_{i \in I} \mathcal{X}_i$

- into a quasi-measurable space by putting: $\left(\prod_{i \in I} \mathcal{X}_i \right)^\Omega := \prod_{i \in I} \mathcal{X}_i^\Omega$

- product random variables on the product are of the form:

- $X(\omega) = (X_i(\omega))_{i \in I}$ with $X_i \in \mathcal{X}_i^\Omega$ for all $i \in I$.

The Function Space

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ and $(\mathcal{Z}, \mathcal{Z}^\Omega)$ two quasi-measurable spaces. We put:
 - $\mathcal{X}^{\mathcal{Z}} := \text{QMS} \left((\mathcal{Z}, \mathcal{Z}^\Omega), (\mathcal{X}, \mathcal{X}^\Omega) \right)$
 - $(\mathcal{X}^{\mathcal{Z}})^\Omega := \left\{ X : \Omega \rightarrow \mathcal{X}^{\mathcal{Z}} \mid ((\omega, z) \mapsto X(\omega)(z)) \in \text{QMS}(\Omega \times \mathcal{Z}, \mathcal{X}) \right\}$
 - function-valued random variables are defined via the product structure
- Then $(\mathcal{X}^{\mathcal{Z}}, (\mathcal{X}^{\mathcal{Z}})^\Omega)$ is a quasi-measurable space.
- Note that such a construction was not possible for measurable spaces!!!

Currying, Uncurrying and the Evaluation Map

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$, $(\mathcal{Y}, \mathcal{Y}^\Omega)$ and $(\mathcal{Z}, \mathcal{Z}^\Omega)$ be quasi-measurable spaces.
- We can then **curry** and **uncurry**:
 - $\text{QMS}(\mathcal{Z} \times \mathcal{Y}, \mathcal{X}) \cong \text{QMS}(\mathcal{Y}, \mathcal{X}^{\mathcal{Z}}) = \text{QMS}\left(\mathcal{Y}, \text{QMS}(\mathcal{Z}, \mathcal{X})\right)$
- In particular, the **evaluation map** is quasi-measurable:
 - $\text{ev} : \mathcal{X}^{\mathcal{Z}} \times \mathcal{Z} \rightarrow \mathcal{X}, \quad \text{ev}(g, z) := g(z).$
- Note that this was not possible in Meas for measurable spaces!!!

More Category-theoretical Constructions

- Similarly, we can define the following in QMS:
 - **coproducts, equalizers, coequalizers**, thus:
 - all small **limits** and all small **colimits**
- even more, we get in QMS:
 - **fibre products**: $\mathcal{X} \times_{\mathcal{S}} \mathcal{Y} \rightarrow \mathcal{S}$,
 - **internal homs**: $\mathcal{Q}_{\mathcal{S}}(\mathcal{X}, \mathcal{Y}) \rightarrow \mathcal{S}$
 - $\text{QMS}_{\mathcal{S}}(\mathcal{X} \times_{\mathcal{S}} \mathcal{Y}, \mathcal{I}) = \text{QMS}_{\mathcal{S}}(\mathcal{X}, \mathcal{Q}_{\mathcal{S}}(\mathcal{Y}, \mathcal{I}))$.
- Note that the latter was not possible in Meas for measurable spaces!!!

Main Theorems

- Theorem: The category of quasi-measurable spaces QMS forms a **quasitopos**, and, is in particular, **locally cartesian closed**.
- Remark: This means that, besides **simply typed λ -calculus**, we get a **dependent type theory** for QMS. Roughly speaking, this means that we can model programs that can vary the output type/space dependent on the input. This makes it easy to implement all result obtained inside QMS in a theorem prover like Lean, Agda or Coq, etc.
- Theorem: The category of quasi-measurable spaces QMS forms a **Heyting category**.
- Remark: QMS has thus an internal logic of a (typed) **intuitionistic first-order logic**.
- Remark: Note that most of this is not true for the category of measurable spaces Meas!!!

Definition - Quasitopos

- A **quasitopos** is a category that:
 - has all finite limits,
 - has all finite colimits,
 - is locally cartesian closed,
 - has a subobject classifier for strong monomorphisms.

The Sample Space - Act 2 - The σ -Algebra

- We now endow the Sample Space (Ω, Ω^Ω) with an additional σ -algebra \mathcal{B}_Ω such that:
 - $\Omega^\Omega \subseteq \text{Meas}((\Omega, \mathcal{B}_\Omega), (\Omega, \mathcal{B}_\Omega))$.
- The **Sample Space** is now the triple: $(\Omega, \Omega^\Omega, \mathcal{B}_\Omega)$.
- Standard example:
 - $\Omega^\Omega = \text{Meas}((\Omega, \mathcal{B}_\Omega), (\Omega, \mathcal{B}_\Omega))$

Topological and Measurable Spaces as Quasi-Measurable Spaces

- If $(\mathcal{X}, \mathcal{E}_{\mathcal{X}})$ is a measurable space or a topological space, etc., then we can turn this into a quasi-measurable space via allowing for the following random variables:
 - $\mathcal{X}^{\Omega} := \mathcal{F}(\mathcal{E}_{\mathcal{X}}) := \{X : \Omega \rightarrow \mathcal{X} \mid \forall A \in \mathcal{E}_{\mathcal{X}}. X^{-1}(A) \in \mathcal{B}_{\Omega}\}$
- Note that the later introduced σ -algebra $\mathcal{B}_{\mathcal{X}}$ might be strictly bigger than the one we started with to turn $(\mathcal{X}, \mathcal{E}_{\mathcal{X}})$ into quasi-measurable space $(\mathcal{X}, \mathcal{X}^{\Omega})$:
 - $\mathcal{E}_{\mathcal{X}} \subsetneq \mathcal{B}_{\mathcal{X}} := \mathcal{B}(\mathcal{X}^{\Omega})$

The σ -Algebra

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ be a quasi-measurable space.
- Then the **induced σ -algebra** is:
 - $\mathcal{B}_\mathcal{X} := \{A \subseteq \mathcal{X} \mid \forall X \in \mathcal{X}^\Omega . X^{-1}(A) \in \mathcal{B}_\Omega\}$
- We can then define the set of admissible random variables with values in $\mathcal{B}_\mathcal{X}$ via:
 - $(\mathcal{B}_\mathcal{X})^\Omega := \{\Psi : \Omega \rightarrow \mathcal{B}_\mathcal{X} \mid \exists D \in \mathcal{B}_{\Omega \times \mathcal{X}} \forall \omega \in \Omega . \Psi(\omega) = D_\omega\} \cong \mathcal{B}_{\Omega \times \mathcal{X}}$
 - where $D_\omega := \{x \in \mathcal{X} \mid (\omega, x) \in D\}$
- Then $(\mathcal{B}_\mathcal{X}, (\mathcal{B}_\mathcal{X})^\Omega)$ is a quasi-measurable space.
- Note that this was not possible in the category of measurable spaces!!!

Theorem - The Adjunction

- A **map** $g : \mathcal{X} \rightarrow \mathcal{Y}$ from a quasi-measurable space $(\mathcal{X}, \mathcal{X}^\Omega)$ to a measurable space $(\mathcal{Y}, \mathcal{B}_\mathcal{Y})$ is
 - **measurable if and only if it is quasi-measurable,**
 - *provided* we use the corresponding choices:
 - $\mathcal{B}_\mathcal{X} := \mathcal{B}(\mathcal{X}^\Omega) := \{A \subseteq \mathcal{X} \mid \forall X \in \mathcal{X}^\Omega . X^{-1}(A) \in \mathcal{B}_\Omega\},$
 - $\mathcal{Y}^\Omega := \mathcal{F}(\mathcal{B}_\mathcal{Y}) := \{Y : \Omega \rightarrow \mathcal{Y} \mid \forall B \in \mathcal{B}_\mathcal{Y} . Y^{-1}(B) \in \mathcal{B}_\Omega\}.$
- In other words, we have the natural identification of sets of maps:
 - $\text{Meas} \left((\mathcal{X}, \mathcal{B}(\mathcal{X}^\Omega)), (\mathcal{Y}, \mathcal{B}_\mathcal{Y}) \right) = \text{QMS} \left((\mathcal{X}, \mathcal{X}^\Omega), (\mathcal{Y}, \mathcal{F}(\mathcal{B}_\mathcal{Y})) \right).$

The Sample Space - Act 3 - Probability Measures

- We now endow the Sample Space $(\Omega, \Omega^\Omega, \mathcal{B}_\Omega)$ with some additional set of **product compatible probability measures** \mathcal{P} on \mathcal{B}_Ω , i.e. such that:
 - for all $P \in \mathcal{P}$ and $D \in \mathcal{B}_{\Omega \times \Omega}$ the map:
 - $\Omega \rightarrow [0,1], \quad \omega \mapsto P(D^\omega),$ is (quasi-)measurable,
 - where $D^\omega := \{\tilde{\omega} \in \Omega \mid (\tilde{\omega}, \omega) \in D\},$
 - for all $P_1, P_2 \in \mathcal{P}$ there exist $\Phi_1, \Phi_2 \in \Omega^\Omega$ and $P \in \mathcal{P}$ such that:
 - $P_1 \otimes P_2 = P(\Phi_1, \Phi_2)$ on $\mathcal{B}_{\Omega \times \Omega},$ i.e. for all $D \in \mathcal{B}_{\Omega \times \Omega}$ we have:
 - $(P_1 \otimes P_2)(D) := \int P_1(D^\omega) P_2(d\omega) = P(\{\omega \in \Omega \mid (\Phi_1(\omega), \Phi_2(\omega)) \in D\}).$
- The **Sample Space** is now the quadruple: $(\Omega, \Omega^\Omega, \mathcal{B}_\Omega, \mathcal{P}).$

The Space of Push-forward Probability Measures

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ be a quasi-measurable space. Define:
 - $\mathcal{P}(\mathcal{X}) := \mathcal{P}(\mathcal{X}, \mathcal{X}^\Omega) := \{P(X) : \mathcal{B}_{\mathcal{X}} \rightarrow [0,1] \mid X \in \mathcal{X}^\Omega, P \in \mathcal{P}\}$
 - $\mathcal{P}(\mathcal{X})^\Omega := \mathcal{P}(\mathcal{X}, \mathcal{X}^\Omega)^\Omega := \{P(X|I) \mid X \in (\mathcal{X}^\Omega)^\Omega, P \in \mathcal{P}\}$
 - $P(X \in A \mid I = \omega) := P\left(\{\tilde{\omega} \in \Omega \mid X(\omega)(\tilde{\omega}) \in A\}\right)$ for $A \in \mathcal{B}_{\mathcal{X}}$
- Lemma: $(\mathcal{P}(\mathcal{X}), \mathcal{P}(\mathcal{X})^\Omega)$ is also a quasi-measurable space.

The Spaces of Markov Kernels and Random Functions

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ and $(\mathcal{Z}, \mathcal{Z}^\Omega)$ be quasi-measurable spaces.
- Then the **space of Markov kernels** from $(\mathcal{Z}, \mathcal{Z}^\Omega)$ to $(\mathcal{X}, \mathcal{X}^\Omega)$:
 - $\mathcal{P}(\mathcal{X})^\mathcal{Z} = \text{QMS} \left((\mathcal{Z}, \mathcal{Z}^\Omega), (\mathcal{P}(\mathcal{X}), \mathcal{P}(\mathcal{X})^\Omega) \right)$
 - is again a quasi-measurable space.
- Also the space of probability distribution over functions:
 - $\mathcal{P}(\mathcal{X}^\mathcal{Z})$ is again a quasi-measurable space.
- Note that these construction were not possible in the category of measurable spaces!!!

Some surprising Lemmata

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ and $(\mathcal{Y}, \mathcal{Y}^\Omega)$ be quasi-measurable spaces.
- Then the following maps are all quasi-measurable:
 - $\mathcal{Y}^{\mathcal{X}} \times \mathcal{B}_{\mathcal{Y}} \rightarrow \mathcal{B}_{\mathcal{X}}, \quad (f, B) \mapsto f^{-1}(B).$
 - $\mathcal{P}(\mathcal{X}) \times \mathcal{B}_{\mathcal{X}} \rightarrow [0,1], \quad (P, A) \mapsto P(A).$
 - $\mathcal{Y}^{\mathcal{X}} \times \mathcal{P}(\mathcal{X}) \rightarrow \mathcal{P}(\mathcal{Y}), \quad (f, P) \mapsto f_*P.$
 - $[0,\infty]^{\mathcal{X}} \times \mathcal{P}(\mathcal{X}) \rightarrow [0,\infty], \quad (h, P) \mapsto \int h(x) P(dx).$
- Note that such statements were not known or even possible in the category of measurable spaces!!!

Theorem: The Product of Markov Kernels

- Assume that there exists an isomorphism of quasi-measurable spaces:
 - $\Omega \times \Omega \cong \Omega$.
- Then for all quasi-measurable spaces $(\mathcal{X}, \mathcal{X}^\Omega)$, $(\mathcal{Y}, \mathcal{Y}^\Omega)$, $(\mathcal{Z}, \mathcal{Z}^\Omega)$ the **product of Markov kernels**:

$$\bullet \otimes : \mathcal{P}(\mathcal{X})^{\mathcal{Y} \times \mathcal{Z}} \times \mathcal{P}(\mathcal{Y})^{\mathcal{Z}} \rightarrow \mathcal{P}(\mathcal{X} \times \mathcal{Y})^{\mathcal{Z}}$$

$$(P(X|Y, Z) \otimes Q(Y|Z))(D|z) := \int P(X \in D^y | Y = y, Z = z) Q(Y \in dy | Z = z)$$

- is a well-defined quasi-measurable map.

Theorem: Strong Probability Monad

- If $\Omega \times \Omega \cong \Omega$ then the triple $(\mathcal{P}, \delta, \mathbb{M})$ is a **strong probability monad** on the cartesian closed category QMS, where:

- $\delta : \mathcal{X} \rightarrow \mathcal{P}(\mathcal{X}),$ $\delta_x(A) := \mathbb{1}_A(x),$

- $\mathbb{M} : \mathcal{P}(\mathcal{P}(\mathcal{X})) \rightarrow \mathcal{P}(\mathcal{X}),$ $\mathbb{M}(\Pi)(A) := \int P(A) d\Pi(P).$

- This thus allows for a notion of computation of monadic type and simply typed λ -calculus.
- We thus get semantics for higher-order probability theory for probabilistic programming language.

Construction of well-behaved Sample Spaces

- Theorem: Let Ω_0 be a set, and:
 - \mathcal{E}_0 a countable set of subsets of Ω_0 that separates the points of Ω_0 .
 - $\Omega := \prod_{n \in \mathbb{N}} \Omega_0$, and $\mathcal{E} := \{\text{pr}_n^{-1}(A) \mid A \in \mathcal{E}_0, n \in \mathbb{N}\}$,
 - $\tilde{\mathcal{P}} := \{P \text{ complete perfect probability measure on } \Omega, \mathcal{E} \subseteq \mathcal{B}_P\}$,
 - $\mathcal{B}_\Omega := \bigcap_{P \in \tilde{\mathcal{P}}} \mathcal{B}_P$, the perfect-universal completion of \mathcal{E} ,
 - $\Omega^\Omega := \text{Meas}((\Omega, \mathcal{B}_\Omega), (\Omega, \mathcal{B}_\Omega))$, $\mathcal{P} := \tilde{\mathcal{P}}|_{\mathcal{B}_\Omega}$
- Then $(\Omega, \Omega^\Omega, \mathcal{B}_\Omega, \mathcal{P})$ satisfies all points of act 1-3 and $\Omega \times \Omega \cong \Omega$.

Fubini Theorem

- Let $(\Omega, \Omega^\Omega, \mathcal{B}_\Omega, \mathcal{P})$ be the sample space from the last slide.
- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ and $(\mathcal{Y}, \mathcal{Y}^\Omega)$ be quasi-measurable spaces and:
 - $f \in [0, \infty]^{\mathcal{X} \times \mathcal{Y}}$, $P \in \mathcal{P}(\mathcal{X})$ and $Q \in \mathcal{P}(\mathcal{Y})$.
- Then we have the equality:

$$\bullet \int \int f(x, y) P(dx) Q(dy) = \int \int f(x, y) Q(dy) P(dx).$$

The Sample Space - Act 4 - The Universal Hilbert Cube

- $\Omega = [0,1]^{\mathbb{N}} = \prod_{n \in \mathbb{N}} [0,1]$, the **Hilbert Cube**,
- \mathcal{B}_{Ω} = set of all *universally measurable* subsets of Ω .
 - Note that this is bigger than the Borel σ -algebra on Ω .
- \mathcal{P} = all probability measures on \mathcal{B}_{Ω} , $\Omega^{\Omega} = \text{Meas} \left((\Omega, \mathcal{B}_{\Omega}), (\Omega, \mathcal{B}_{\Omega}) \right)$.
- We call this Sample Space $(\Omega, \Omega^{\Omega}, \mathcal{B}_{\Omega}, \mathcal{P})$ the **Universal Hilbert Cube**.
- Interpretation: Countably infinite sequence of uniformly distributed samples (e.g. from a (pseudo-)random number generator).
- Note that it satisfies act 1-3 and the iso: $\Omega \times \Omega \cong \Omega$ (via “Hilbert’s Hotel”).

The Category of Quasi-Universal Spaces

- Definition: A **quasi-universal space** $(\mathcal{X}, \mathcal{X}^\Omega)$ is - per definition - just a quasi-measurable space where the sample space Ω is the **universal Hilbert cube**.
- We abbreviate the category of quasi-universal spaces as QUS.

Countably Separated and Standard Quasi-Measurable Spaces

- Definition: A quasi-measurable space $(\mathcal{X}, \mathcal{X}^\Omega)$ is called:
 - **countably separated** if there exists a countable subset $\mathcal{E} \subseteq \mathcal{B}_\mathcal{X}$ that separates the points of \mathcal{X} .
 - **standard quasi-measurable space** if there are quasi-measurable maps:
 - $\iota : (\mathcal{X}, \mathcal{X}^\Omega) \rightarrow (\Omega, \Omega^\Omega)$ and $r : (\Omega, \Omega^\Omega) \rightarrow (\mathcal{X}, \mathcal{X}^\Omega)$ s.t.:
 - $r \circ \iota = \text{id}_\mathcal{X}$.

Theorem: Disintegration of Markov Kernels

- Let $(\mathcal{X}, \mathcal{X}^\Omega)$ and $(\mathcal{Y}, \mathcal{Y}^\Omega)$ and $(\mathcal{Z}, \mathcal{Z}^\Omega)$ be quasi-universal spaces.
 - Let $(\mathcal{Y}, \mathcal{Y}^\Omega)$ be *countably separated*. and:
 - either $(\mathcal{X}, \mathcal{X}^\Omega)$ or $(\mathcal{Z}, \mathcal{Z}^\Omega)$ be a *standard* quasi-universal space.
- Then the product of Markov kernels:
 - $\otimes : \mathcal{P}(\mathcal{X})^{\mathcal{Y} \times \mathcal{Z}} \times \mathcal{P}(\mathcal{Y})^{\mathcal{Z}} \rightarrow \mathcal{P}(\mathcal{X} \times \mathcal{Y})^{\mathcal{Z}}$
 - is a (surjective) quotient map of quasi-universal spaces.
- More concretely, for every $P(X, Y | Z) \in \mathcal{P}(\mathcal{X} \times \mathcal{Y})^{\mathcal{Z}}$ there exists $P(X | Y, Z) \in \mathcal{P}(\mathcal{X})^{\mathcal{Y} \times \mathcal{Z}}$ such that: $P(X, Y | Z) = P(X | Y, Z) \otimes P(Y | Z)$.

Conditional Kolmogorov Extension Theorem

- Let $(\mathcal{X}_n, \mathcal{X}_n^\Omega)$, $n \in \mathbb{N}$, a sequence of *standard* quasi-universal spaces and $(\mathcal{Z}, \mathcal{Z}^\Omega)$ be any quasi-universal space.
- Assume we have $Q_n(X_{0:n} | Z) \in \mathcal{P}(\mathcal{X}_{0:n})^{\mathcal{Z}}$ such that for every $n \in \mathbb{N}$:
 - $\text{pr}_{0:n,*} Q_{n+1}(X_{0:n+1} | Z) = Q_n(X_{0:n} | Z)$.
- Then there exists a unique $Q(X_{\mathbb{N}} | Z) \in \mathcal{P}(\mathcal{X}_{\mathbb{N}})^{\mathcal{Z}}$ such that:
 - $\text{pr}_{0:n,*} Q(X_{0:n+1} | Z) = Q_n(X_{0:n} | Z)$ for all $n \in \mathbb{N}$,
- where $\mathcal{X}_{\mathbb{N}} := \prod_{n \in \mathbb{N}} \mathcal{X}_n$.

Conditional De Finetti Theorem

- $(\mathcal{X}, \mathcal{X}^\Omega)$ standard quasi-universal spaces, $(\mathcal{Z}, \mathcal{Z}^\Omega)$ any quasi-universal space.
- For a Markov kernel $Q(X_{\mathbb{N}} | Z) \in \mathcal{P}(\mathcal{X}^{\mathbb{N}})^{\mathcal{Z}}$ the following is equivalent:
 - $Q(X_{\mathbb{N}} | Z)$ is **exchangable**, i.e. invariant under all finite permutations: $\rho : \mathbb{N} \cong \mathbb{N}$.
 - There exists a quasi-universal space \mathcal{Y} and $K(X | Y) \in \mathcal{P}(\mathcal{X})^{\mathcal{Y}}$ and $P(Y | Z) \in \mathcal{P}(\mathcal{Y})^{\mathcal{Z}}$ such that :

$$\bullet Q(X_{\mathbb{N}} | Z) = \left(\bigotimes_{n \in \mathbb{N}} K(X_n | Y) \right) \circ P(Y | Z).$$

- In this case we can w.l.o.g. take: $\mathcal{Y} = \mathcal{P}(\mathcal{X})$ and $K(X \in A | Y = P) := P(A)$.

Transitional Conditional Independence

- Consider a Markov kernel: $P(X, Y, Z | T) \in \mathcal{P}(\mathcal{X} \times \mathcal{Y} \times \mathcal{Z})^{\mathcal{T}}$.
- We say that X is **conditional independent** of Y given Z w.r.t. $P(X, Y, Z | T)$,
 - in symbols: $X \perp\!\!\!\perp Y | Z$ if:
 - there *exists* a Markov kernel $Q(X | Z) \in \mathcal{P}(\mathcal{X})^{\mathcal{Z}}$ such that:
 - $P(X, Y, Z | T) = Q(X | Z) \otimes P(Y, Z | T)$.

Partially Generic Causal Bayesian Networks

- A **partially generic causal Bayesian network** - per definition - consists of:
 - a **conditional directed acyclic graph (CDAG)**: $G = (J, V, E)$,
 - an input variable X_j on a quasi-universal space \mathcal{X}_j for each $j \in J$,
 - an output variable X_v on a *standard* quasi-universal space \mathcal{X}_v for each $v \in V$,
 - an **exceptional set**: $W \subseteq V$,
 - a Markov kernel: $P_v(X_v | X_{\text{Pa}^G(v)}) \in \mathcal{P}(\mathcal{X}_v)^{\mathcal{X}_{\text{Pa}^G(v)}}$ for $v \in V \setminus W$.

Partially Generic Causal Bayesian Networks

- For a partially generic causal Bayesian network with exceptional set W we introduce for $w \in W$:
 - an **indicator variable**: $I_w \rightarrow w$,
 - a quasi-universal space: $\mathcal{X}_{I_w} := \mathcal{P}(\mathcal{X}_w)^{\mathcal{X}_{\text{Pa}G(w)}}$,
 - a “**generic**” Markov kernel:
 - $P_w \left(X_w \in A \mid X_{\text{Pa}G(w)} = x, X_{I_w} = Q \right) := Q \left(X_w \in A \mid X_{\text{Pa}G(w)} = x \right)$.
- So we get a joint Markov kernel: $P(X_V, X_J, X_{I_W} \mid X_J, X_{I_W})$.

Theorem: Global Markov Property

- For every partially generic causal Bayesian network with exceptional set W and any subsets: $A, B, C \subseteq V \cup I_W \cup J$ we have the implication:

- $A \perp B | C \implies X_A \perp\!\!\!\perp X_B | X_C.$

**(Proposed)
Answers**

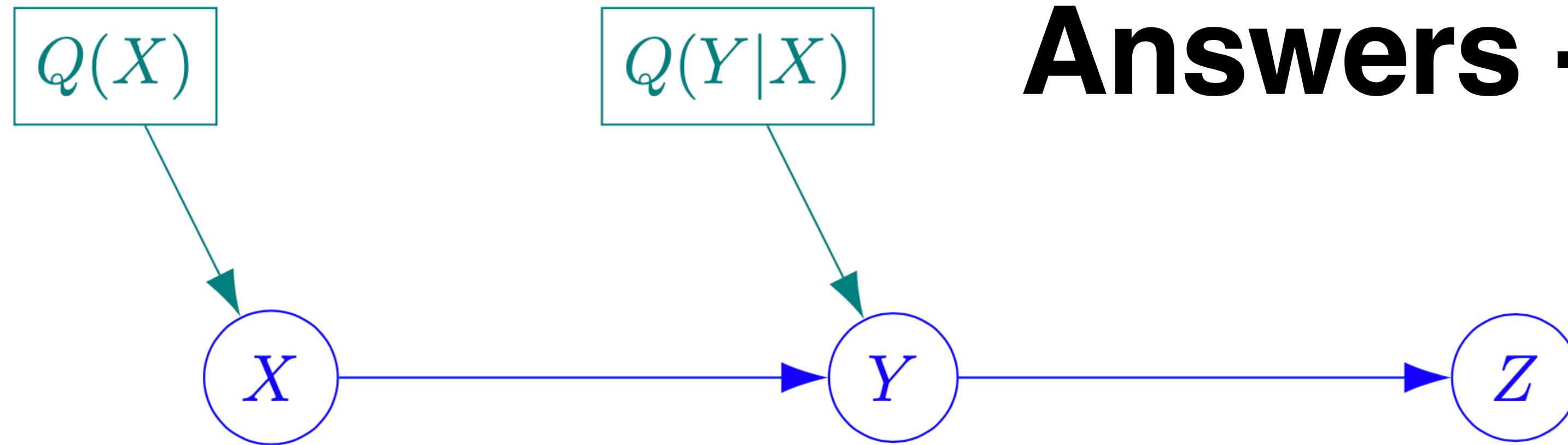
Answers - Stochastic Process

- Definition: A **stochastic process** is a quasi-measurable map:
 - $X : \Omega \rightarrow \mathcal{X}^{\mathcal{T}}, \quad \omega \mapsto (t \mapsto X(\omega)(t)).$
- Lemma: This is *equivalent* to a quasi-measurable map: $X : \Omega \times \mathcal{T} \rightarrow \mathcal{X}, \quad (\omega, t) \mapsto X(\omega, t).$
- Lemma: The map: $\mathcal{X}^{\mathcal{T}} \rightarrow \prod_{t \in \mathcal{T}} \mathcal{X}, \quad X \mapsto (X(t))_{t \in \mathcal{T}},$ is quasi-measurable.
- Lemma: If $T : \Omega \rightarrow \mathcal{T}$ is quasi-measurable (random time) then the map:
 - $\Omega \rightarrow \mathcal{X}, \quad \omega \mapsto X(\omega)(T(\omega))$ is again quasi-measurable.

Answers - Probabilistic Programs

- Definition: A **probabilistic program** with input $x \in \mathcal{X}$ and output $z \in \mathcal{Z}$ is quasi-measurable map: $\mathcal{X} \rightarrow \mathcal{P}(\mathcal{Z})$.
- Theorem: We have the natural **curry / uncurry isomorphism**:
 - $\text{QMS}(\mathcal{X} \times \mathcal{Y}, \mathcal{P}(\mathcal{Z})) \cong \text{QMS}(\mathcal{X}, \text{QMS}(\mathcal{Y}, \mathcal{P}(\mathcal{Z})))$
- Theorem: QMS is a **quasitopos**, thus allows for **dependent type theory**.
- Theorem: The triple $(\mathcal{P}, \delta, \mathbb{M})$ forms a **strong probability monad** on the category of quasi-measurable spaces QMS (for certain sample spaces, e.g. the universal Hilbert cube). Thus allows for **higher-order probabilistic programs**.

Answers - Graphical Models



- **Partially generic causal Bayesian networks** can model graphical models with non-random input variables.
- **Transitional conditional independence** also works with non-random input variables.
- Theorem: **Global Markov Property**: For $A, B, C \subseteq V \cup I_W \cup J$ we have:
 - $A \perp B | C \implies X_A \perp\!\!\!\perp X_B | X_C$.
- Example: Here $Q(Y|X)$ is a non-random input variable with values in $\mathcal{L} := \text{QUS}(\mathcal{X}, \mathcal{P}(\mathcal{Y}))$
 - Then Y is determined by the new quasi-measurable mechanism:
 - $\mathcal{L} \times \mathcal{X} \rightarrow \mathcal{P}(\mathcal{Y}), \quad (Q(Y|X), x) \mapsto Q(Y|X = x)$.
 - We can now read off the graph: $Z \perp\!\!\!\perp X, Q(Y|X), Q(X) | Y$.

Answers - Causal Assumptions

- Model **potential outcome** as quasi-measurable map / random function:
 - $G : \Omega \rightarrow \mathcal{Y}^{\mathcal{X}}$
- Potential outcome under treatment $X = x$ then: $Y_x := G(x)$.
- Rephrase causal assumptions:
 - Strong Ignorability: $X \perp\!\!\!\perp G \mid Z$,
 - Consistency: $Y = G(X)$.
- Everything is well-defined and quasi-measurable.

Answers - Counterfactual Probabilities

- Theorem: Disintegration of Markov kernels.
- Model potential outcome as: $G \in (\mathcal{Y}^{\mathcal{X}})^{\Omega}$
- Assume that \mathcal{X} to countably separated quasi-universal space.
- Then via the disintegration theorem there exists conditional:
 - $P(G | X) \in \mathcal{P}(\mathcal{G})^{\mathcal{X}}$ such that $P(G, X) = P(G | X) \otimes P(X)$.
- Evaluation maps and push-forwards are quasi-measurable, which implies:
 - $C(A | x, x') := P(G(x) \in A | X = x')$ defines:
 - well-defined and quasi-measurable $C \in \mathcal{P}(\mathcal{Y})^{\mathcal{X} \times \mathcal{X}}$
- So, conditional counterfactual probabilities are well-defined and quasi-measurable.

Answers - Statistics and Probability Theory

- For (standard) quasi-universal spaces we at least can do the following:
 - Theorem: Disintegration of Markov kernels.
 - Remark: This allows for Bayes' Rule and thus Bayesian Statistics.
 - Theorem: Fubini Theorem.
 - Theorem: Conditional de Finetti Theorem.
 - Theorem: Kolmogorov Extension Theorem.
 - Theorem: Global Markov Property for graphical models like partially generic causal Bayesian networks.

Recommendation

- For probabilistic programming, graphical models, causality, statistics, etc.
 - use for:
 - sample space \rightarrow the universal Hilbert cube
 - replace:
 - measurable spaces \rightarrow quasi-measurable spaces
 - measurable maps \rightarrow quasi-measurable maps
 - categorical construction in \mathbf{Meas} \rightarrow categorical construction in \mathbf{QMS}
- study more of the (classical) theory in this framework (e.g. martingales).

• Patrick Forré, *Quasi-Measurable Spaces*, 2021, <https://arxiv.org/abs/2109.11631>.

More about Convenient Categories

- Probability Theory

- **Quasi-Borel Spaces** - by Chris Heunen, Ohad Kammar, Sam Staton, Hongseok Yang
- **Quasi-Measurable Spaces** - by Patrick Forré, <https://arxiv.org/abs/2109.11631>

- Topology

- **Compactly Generated Weakly Hausdorff Spaces (CGWH)** - by Witold Hurewicz, David Gale, Norman Steenrod, John C. Moore, Michael C. McCord, Neil Strickland, et al ([script](#))
- **Condensed Sets** - by Peter Scholz, Dustin Clausen ([script](#))

- Differential Geometry

- **Diffeological Spaces** - by Kuo Tsai Chen, Jean-Marie Souriau, Patrick Iglesias-Zemmour, John Baez, Alexander Hoffnung, Andrew Stacey, et al.

Thank you for your attention!